

Diagrammatics for F_4

$$\begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ \diagup \diagdown \end{array} + \begin{array}{c} \diagup \diagdown \\ \diagup \diagdown \end{array} = \frac{2\alpha}{\delta + 2} \left(\begin{array}{c} | \\ | \end{array} + \begin{array}{c} \cup \\ \cap \end{array} + \begin{array}{c} \times \\ \times \end{array} \right)$$

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Outline

Goal: Define an F_4 analogue of the oriented and unoriented Brauer categories and algebras.

Overview:

- 1 Centralizers in classical types
- 2 Diagrammatic categories
- 3 Oriented and unoriented Brauer categories
- 4 The category \mathcal{F}
- 5 The functor $\mathcal{F} \rightarrow \mathfrak{f}\text{-mod}$
- 6 Further directions

Centralizers in type A

For $\delta \in \mathbb{N}$, let $V = \mathbb{C}^\delta$ be the natural \mathfrak{gl}_δ -module.

The natural permutation action yields a surjective algebra homomorphism

$$\mathbb{C}\mathfrak{S}_r \twoheadrightarrow \text{End}_{\mathfrak{gl}_\delta}(V^{\otimes r}).$$

This is part of **Schur–Weyl duality**.

If we want to include the dual module V^* , we have

$$\text{OB}_{r,s}(\delta) \twoheadrightarrow \text{End}_{\mathfrak{gl}_\delta}(V^{\otimes r} \otimes (V^*)^{\otimes s}),$$

where $\text{OB}_{r,s}(\delta)$ is the **walled Brauer algebra** (or **oriented Brauer algebra**).

Centralizers in other types

Types BCD

Let G be an orthogonal or symplectic Lie group.

The natural module V is **self-dual**. Let $\delta = \dim V$.

We have

$$B_r(\delta) \twoheadrightarrow \text{End}_G(V^{\otimes r}),$$

where $B_r(\delta)$ is the **Brauer algebra**.

Exceptional types

Less is known in exceptional types.

We will focus on the Lie algebra \mathfrak{f} of **type F_4** .

If V is the natural \mathfrak{f} -module, we want

$$?? \twoheadrightarrow \text{End}_{\mathfrak{f}}(V^{\otimes r})$$

Bigger picture

In fact, things become much more natural if we describe **more**.

General idea

Instead of only considering endomorphisms of $V^{\otimes r}$, we consider **all** morphisms $V^{\otimes r} \rightarrow V^{\otimes s}$ (and include dual V^* in type A).

Want a “nice” **monoidal category** \mathcal{C} , together with a full and essentially surjective monoidal functor

$$\mathcal{C} \rightarrow \mathfrak{g}\text{-mod.}$$

Desired features

- 1 \mathcal{C} should be easy to describe, e.g. want a **presentation** in terms of generators and relations.
- 2 Should recover above centralizer statements when restricting to appropriate morphism spaces.

Strict monoidal categories

A **strict monoidal category** is a category \mathcal{C} equipped with

- a bifunctor (the **tensor product**) $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and
- a **unit object** $\mathbb{1}$,

such that, for objects A, B, C and morphisms f, g, h ,

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$,
- $(f \otimes g) \otimes h = f \otimes (g \otimes h)$,
- $1_{\mathbb{1}} \otimes f = f = f \otimes 1_{\mathbb{1}}$.

Remark: Non-strict monoidal categories

In a (not necessarily strict) **monoidal category**, the equalities above are replaced by isomorphism, and we impose some **coherence conditions**.

Every monoidal category is monoidally equivalent to a strict one.

Linear monoidal categories

For simplicity, we work over the ground field \mathbb{C} .

A **strict linear monoidal category** is a strict monoidal category such that

- each morphism space is a \mathbb{C} -module,
- composition of morphisms is \mathbb{C} -bilinear,
- tensor product of morphisms is \mathbb{C} -bilinear.

The interchange law

The axioms of a strict monoidal category imply the **interchange law**: For $A_1 \xrightarrow{f} A_2$ and $B_1 \xrightarrow{g} B_2$, the following diagram commutes:

$$\begin{array}{ccc} A_1 \otimes B_1 & \xrightarrow{1 \otimes g} & A_1 \otimes B_2 \\ f \otimes 1 \downarrow & \searrow f \otimes g & \downarrow f \otimes 1 \\ A_2 \otimes B_1 & \xrightarrow{1 \otimes g} & A_2 \otimes B_2 \end{array}$$

String diagrams

Fix a strict monoidal category \mathcal{C} .

We will denote a morphism $f: A \rightarrow B$ by:



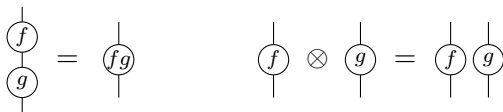
The **identity map** $1_A: A \rightarrow A$ is a string with no label:



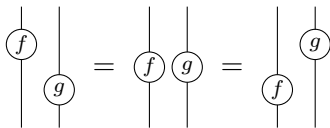
We sometimes omit the object labels when they are clear or unimportant.

String diagrams

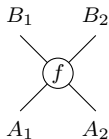
Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:



The **interchange law** then becomes:



A morphism $f: A_1 \otimes A_2 \rightarrow B_1 \otimes B_2$ can be depicted:



Presentations of strict monoidal categories

One can give **presentations** of some strict \mathbb{C} -linear monoidal categories, just as for monoids, groups, algebras, etc.

Objects: If the objects are generated by some collection $A_i, i \in I$, then we have all possible tensor products of these objects:

$$\mathbb{1}, \quad A_i, \quad A_i \otimes A_j \otimes A_k \otimes A_\ell, \quad \text{etc.}$$

Morphisms: If the morphisms are generated by some collection $f_j, j \in J$, then we have all possible compositions and tensor products of these morphisms (whenever these make sense):

$$1_{A_i}, \quad f_j \otimes (f_i f_k) \otimes (f_\ell), \quad \text{etc.}$$

We then often impose some **relations** on these morphism spaces.

String diagrams: We can build complex diagrams out of our simple generating diagrams.

Example: monoidally generated symmetric groups

Define a strict monoidal category Sym with one generating object I and denote

$$1_I = |$$

We have one generating morphism

$$\times : I \otimes I \rightarrow I \otimes I.$$

We impose the relations:

$$\begin{array}{c} \text{Diagram 1: A crossing of two strands, with the top-left and bottom-right strands crossing over the top-right and bottom-left strands.} \\ \text{Diagram 2: Two parallel vertical strands.} \end{array} = \quad , \quad \begin{array}{c} \text{Diagram 3: A crossing of two strands, with the top-right and bottom-left strands crossing over the top-left and bottom-right strands.} \\ \text{Diagram 4: A crossing of two strands, with the top-left and bottom-right strands crossing over the top-right and bottom-left strands.} \end{array}.$$

Then

$$\text{End}_{Sym}(I^{\otimes n}) = S_n$$

is the **symmetric group** on n letters.

Example: monoidally generated symmetric groups

This monoidal presentation of S_n is very efficient! We only needed

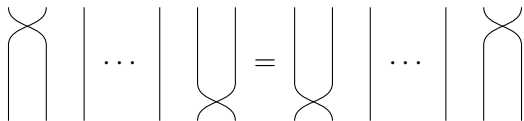
- one generating morphism, and
- two relations,

to get **all** the symmetric groups.

Note that the “distant braid relation”

$$s_i s_j = s_j s_i, \quad |i - j| > 1$$

for simple transpositions follows for free from the interchange law:



Note: If we define Sym to be linear, then $\text{End}_{Sym}(\uparrow^{\otimes n}) = \mathbb{C}S_n$.

Oriented Brauer category

For $\delta \in \mathbb{C}$, the **oriented Brauer category** $\mathcal{OB}(\delta)$ is the linear monoidal category defined as follows.

Two generating objects: \uparrow and \downarrow

Four generating morphisms: \cup , \cap , \nearrow , \searrow

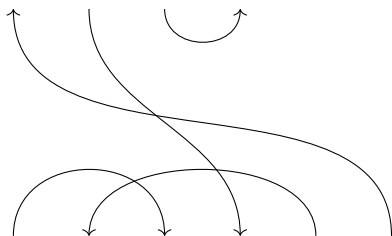
Define $\nearrow := \text{cup} \cap$, $\searrow := \text{cup} \cup$, $\cup := \delta$, $\cap := \delta$.

Relations:

$$\begin{aligned} \text{cup} \cup &= \uparrow \uparrow, & \text{cup} \cap &= \text{cup} \cap, & \text{cup} \cup &= \downarrow \uparrow, & \text{cup} \cap &= \uparrow \downarrow, \\ \text{cup} \cup &= \uparrow, & \text{cup} \cup &= \downarrow, & \text{cup} &= \delta \mathbf{1}_1. \end{aligned}$$

Oriented Brauer category

An arbitrary morphism in $OB(\delta)$ is a \mathbb{C} -linear combination of **oriented Brauer diagrams**. E.g.



$$: \uparrow \otimes \downarrow^{\otimes 3} \otimes \uparrow^{\otimes 2} \rightarrow \uparrow \otimes \downarrow^{\otimes 2} \otimes \uparrow .$$

Composition: vertical “gluing”, replace closed components by a factor of δ .

Tensor product: horizontal concatenation.

Oriented Brauer category

Universal property

Any linear symmetric monoidal category \mathcal{C} with dual objects V, V^* of dimension δ admits a linear monoidal functor

$$OB(\delta) \rightarrow \mathcal{C}, \quad \uparrow \mapsto V, \quad \downarrow \mapsto V^*.$$

Corollary

We have a linear monoidal functor ($V =$ natural module)

$$OB(\delta) \rightarrow \mathfrak{gl}_\delta\text{-mod}, \quad \uparrow \mapsto V, \quad \downarrow \mapsto V^*.$$

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} \mapsto (V^{\otimes 2} \rightarrow V^{\otimes 2}, \quad v \otimes w \mapsto w \otimes v),$$

$$\begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} \mapsto (V^* \otimes V \rightarrow V \otimes V^*, \quad f \otimes v \mapsto v \otimes f),$$

$$\cap \mapsto (V \otimes V^* \rightarrow \mathbb{C}, \quad v \otimes f \mapsto f(v)),$$

$$\cup \mapsto \left(\mathbb{C} \rightarrow V^* \otimes V, \quad 1 \mapsto \sum_{v \in \mathbf{B}_V} v^* \otimes v \right),$$

where \mathbf{B}_V is a basis for V and $v^*, v \in \mathbf{B}_V$, is the dual basis.

Oriented Brauer category

In fact, this functor is **full**. In particular, we recover the surjection

$$\mathrm{OB}_{r,s}(\delta) \cong \mathrm{End}_{\mathcal{O}\mathcal{B}(\delta)}(\uparrow^{\otimes r} \otimes \downarrow^{\otimes s}) \twoheadrightarrow \mathrm{End}_{\mathfrak{gl}_\delta}(V^{\otimes r} \otimes (V^*)^{\otimes s}),$$

where $\mathrm{OB}_{r,s}(\delta)$ is the **walled Brauer algebra** (or **oriented Brauer algebra**).

Fact: Every f.d. \mathfrak{gl}_δ -module is a summand of $V^{\otimes r} \otimes (V^*)^{\otimes s}$ for some r, s .

Philosophy

All homomorphisms between f.d. \mathfrak{gl}_δ -modules are built from the “flip” symmetry and the natural pairing of V and V^* .

Brauer category

Fix $\delta \in \mathbb{C}$. The Brauer category $\mathcal{B}(\delta)$ is the linear monoidal category defined as follows.

One generating object: $\mathbb{1}$

Three generating morphisms: \cup, \cap, \times

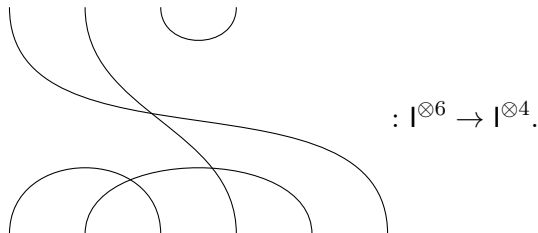
Relations:

$$\begin{aligned} \text{Cup} &= \text{Cup}, & \text{Cross} &= \text{Cross}, & \text{Cap} &= \text{Cap}, & \text{Cup} &= \text{Cup}, \\ \text{Cap} &= \text{Cap}, & \cup &= \delta \mathbb{1}, & \cap &= \delta \mathbb{1} \end{aligned}$$

Brauer category

An arbitrary morphism in $\mathcal{B}(\delta)$ is a linear combination of **Brauer diagrams**.

E.g.



Composition: vertical “gluing”, replace closed components by a factor of δ .

Tensor product: horizontal concatenation.

Brauer category

Universal property

Any linear symmetric monoidal category \mathcal{C} with self-dual object V of dimension δ admits a linear monoidal functor

$$\mathcal{B}(\delta) \rightarrow \mathcal{C}, \quad 1 \mapsto V.$$

Corollary

We have a linear monoidal functor ($V =$ natural module)

$$\mathcal{B}(\delta) \rightarrow \mathfrak{so}_\delta\text{-mod}, \quad 1 \mapsto V.$$

$$\begin{aligned} \times &\mapsto (V^{\otimes 2} \rightarrow V^{\otimes 2}, \quad v \otimes w \mapsto w \otimes v), \\ \cap &\mapsto (V^{\otimes 2} \rightarrow \mathbb{C}, \quad v \otimes w \mapsto \langle v, w \rangle), \\ \cup &\mapsto (\mathbb{C} \rightarrow V^{\otimes 2}, \quad 1 \mapsto \sum_{v \in \mathbf{B}_V} v \otimes v^\vee), \end{aligned}$$

where $\langle \cdot, \cdot \rangle$ is the bilinear form, and $^\vee$ denotes the dual basis.

Brauer category

In fact, this functor is **full**. In particular, we recover the surjection

$$B_r(\delta) \cong \text{End}_{\mathcal{B}(\delta)}(I^{\otimes r}) \twoheadrightarrow \text{End}_{\mathfrak{so}_\delta}(V^{\otimes r}),$$

where $B_r(\delta)$ is the **Brauer algebra**

Fact: Every f.d. \mathfrak{so}_δ -module is a summand of $V^{\otimes r}$ for some r .

Philosophy

All homomorphisms between f.d. \mathfrak{so}_δ -modules are built from the “flip” symmetry and the bilinear pairing.

Similar statements are true in the **symplectic** case.

Goal

We want analogous results in type F_4 .

Thus, we want a linear monoidal category \mathcal{F} and a functor

$$\Phi: \mathcal{F} \rightarrow \mathfrak{f}\text{-mod}$$

such that

- \mathcal{F} has a simple presentation,
- Φ is **full** (surjective on morphisms)
- Φ is **essentially surjective** (surjective on isomorphism classes) after adding in summands.

A preliminary category

For $\alpha, \delta \in \mathbb{C}$, $\mathcal{T} = \mathcal{T}_{\alpha, \delta}$ is the linear monoidal category defined as follows.

One generating object: $\mathbb{1}$ Four generating morphisms: $\cup, \cap, \times, \circ$

Relations:

$$\begin{array}{l}
 \cup = | = \cap, \quad \cup := \cup = \cup, \quad \cap = \cap, \\
 \times = ||, \quad \times = \times, \quad \cup \cap = \cup \cap, \quad \times \times = \times \times, \\
 \cup = \cap, \quad \cup = \cup, \quad \cup = \alpha |, \quad \circ = \delta \mathbb{1}, \quad \cup = 0.
 \end{array}$$

Notes

- Can scale \cup by $\alpha^{-1/2}$ to reduce to case $\alpha = 1$.
- Category is strict pivotal (isotopy invariance), so we can define

$$\cup \cap := \cup \cap = \cup \cap.$$

The category \mathcal{F}

Assume $\delta \neq -2$. Then $\mathcal{F} = \mathcal{F}_{\alpha, \delta}$ is obtained from $\mathcal{T}_{\alpha, \delta}$ by imposing three additional relations:

$$\begin{array}{c} \diagup \diagdown + \diagdown \diagup + \diagup \diagup \diagdown \diagdown = \frac{2\alpha}{\delta+2} \left(\parallel + \cup + \times \right), \end{array}$$

$$\begin{array}{c} \square = \frac{\alpha^2(\delta+14)}{2(\delta+2)^2} \left(\parallel + \cup \right) + \frac{\alpha(\delta-6)}{2(\delta+2)} \left(\diagup \diagdown + \diagdown \diagup \right) + \frac{3\alpha^2(2-\delta)}{2(\delta+2)^2} \times, \end{array}$$

$$\begin{array}{c} \begin{array}{c} \diagup \diagdown \\ | \\ \diagdown \diagup \end{array} = \frac{\alpha(10-\delta)}{4(\delta+2)} \left(\begin{array}{c} \diagup \diagdown \\ | \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagup \diagdown \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagup \diagdown \\ | \\ \diagup \diagdown \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} \right) \\ - \frac{\alpha^2(\delta+30)}{8(\delta+2)^2} \left(\begin{array}{c} \diagdown \\ | \\ \cup \end{array} + \begin{array}{c} \diagup \\ | \\ \diagdown \end{array} + \begin{array}{c} \diagdown \\ | \\ \diagup \end{array} + \begin{array}{c} \cup \\ | \\ \diagdown \end{array} + \begin{array}{c} \cup \\ | \\ \diagup \end{array} \right) \\ + \frac{3\alpha^2(\delta-2)}{8(\delta+2)^2} \left(\begin{array}{c} \cup \\ | \\ \diagdown \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ | \\ \diagdown \diagup \end{array} \right). \end{array}$$

The Albert algebra

Let

$$A_{\mathbb{R}} = \left\{ \begin{pmatrix} \lambda_1 & x_3 & \bar{x}_2 \\ \bar{x}_3 & \lambda_2 & x_1 \\ x_2 & \bar{x}_1 & \lambda_3 \end{pmatrix} : \lambda_i \in \mathbb{R}, x_i \in \mathbb{O} \right\}$$

denote the set of 3×3 self-adjoint matrices over the octonions \mathbb{O} , equipped with the bilinear operation

$$a \circ b := \frac{1}{2}(ab + ba), \quad a, b \in A_{\mathbb{R}},$$

where the juxtaposition ab denotes usual matrix multiplication.

We have a **trace map**

$$\text{tr}: A_{\mathbb{R}} \rightarrow \mathbb{R}, \quad \text{tr} \begin{pmatrix} \lambda_1 & x_3 & \bar{x}_2 \\ \bar{x}_3 & \lambda_2 & x_1 \\ x_2 & \bar{x}_1 & \lambda_3 \end{pmatrix} = \lambda_1 + \lambda_2 + \lambda_3.$$

$A := \mathbb{C} \otimes_{\mathbb{R}} A_{\mathbb{R}}$ is the unique **simple exceptional Jordan algebra**.

The Lie algebra of type F_4

The group G of algebra automorphisms of $A_{\mathbb{R}}$ is the compact connected real Lie group of type F_4 .

Let \mathfrak{f} be the complexification of the Lie algebra of G .

The **symmetric bilinear form**

$$B: A \otimes A \rightarrow \mathbb{C}, \quad B(a \otimes b) := \text{tr}(a \circ b)$$

is nondegenerate and G -invariant (hence \mathfrak{f} -invariant).

Thus, we have a decomposition of \mathfrak{f} -modules

$$A = \mathbb{C}1_A \oplus V, \quad V = \ker(\text{tr}).$$

V is the **natural** \mathfrak{f} -module, and

$$\dim V = 26.$$

The functor

Fix a basis \mathbf{B}_V of V , with dual basis $\{b^\vee : b \in B_V\}$, and let

$$\pi: A = \mathbb{C}1_A \oplus V \rightarrow V,$$

be the natural projection.

Theorem (Gandhi–S.–Zaynullin 2021)

There is a unique monoidal functor

$$\Phi: \mathcal{T}_{7/3,26} \rightarrow \mathfrak{f}\text{-mod}$$

given on objects by $I \mapsto V$ and on morphisms by

$$\Phi(\frown): V \otimes V \rightarrow V, \quad a \otimes b \mapsto \pi(a \circ b),$$

$$\Phi(\bowtie): V \otimes V \rightarrow V \otimes V, \quad a \otimes b \mapsto b \otimes a,$$

$$\Phi(\cup): \mathbb{C} \rightarrow V \otimes V, \quad 1 \mapsto \sum_{b \in \mathbf{B}_V} b \otimes b^\vee,$$

$$\Phi(\cap): V \otimes V \rightarrow \mathbb{C}, \quad a \otimes b \mapsto B(a \otimes b) = \text{tr}(a \circ b).$$

The functor

Theorem (Gandhi–S.–Zaynullin 2021)

- 1 The functor Φ is **full** (surjective on morphisms).
- 2 The functor Φ **factors through** $\mathcal{F}_{7/3,26}$, giving

$$\Phi: \mathcal{F}_{7/3,26} \rightarrow \mathfrak{f}\text{-mod.}$$

Remark

The relation

$$\begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} + \begin{array}{c} \diagdown \diagup \\ \diagup \diagdown \end{array} + \begin{array}{c} | \\ | \\ | \end{array} = \frac{2\alpha}{\delta + 2} \left(\begin{array}{c} | \\ | \\ | \end{array} + \begin{array}{c} \cup \\ \cup \\ \cup \end{array} + \begin{array}{c} \times \\ \times \\ \times \end{array} \right)$$

corresponds to the **Cayley–Hamilton equation** for traceless 3×3 octonionic matrices.

Idempotent completion

Additive envelope

The **additive envelope** $\text{Add}(\mathcal{F})$ is a linear monoidal category with:

Objects: Formal finite direct sums $\bigoplus_{i=1}^n X_i$, $X_i \in \mathcal{F}$.

Morphisms: Morphisms

$$\bigoplus_{i=1}^n X_i \rightarrow \bigoplus_{j=1}^m Y_j$$

are $m \times n$ matrices, where the (j, i) -entry is a morphism $f_{ij}: X_i \rightarrow Y_j$ in \mathcal{F} . Composition is matrix multiplication.

Karoubi envelope

The **additive Karoubi envelope** $\text{Kar}(\mathcal{F})$ is a linear monoidal category with:

Objects: Pairs (X, e) with $X \in \text{Add}(\mathcal{F})$ and $e: X \rightarrow X$ an idempotent.

Morphisms: $\text{Hom}_{\text{Kar}(\mathcal{F})}((X, e), (X', e')) = e' \text{Hom}_{\text{Add}(\mathcal{F})}(X, X')e$.

Idempotent completion

Since the category $\mathfrak{f}\text{-mod}$ is **idempotent complete** (i.e. it contains the images of idempotents), we have an induced functor

$$\text{Kar}(\Phi): \text{Kar}(\mathcal{F}_{7/3,26}) \rightarrow \mathfrak{f}\text{-mod}.$$

Proposition (Gandhi–S.–Zaynullin 2021)

The functor $\text{Kar}(\Phi)$ is **full** and **essentially surjective** (i.e. every finite-dimensional \mathfrak{f} -module is isomorphic to an object in the image of $\text{Kar}(\Phi)$).

Corollary (Centralizer property)

We have a surjective algebra isomorphism

$$\text{End}_{\mathcal{F}_{7/3,26}}(I^{\otimes r}) \twoheadrightarrow \text{End}_{\mathfrak{f}}(V^{\otimes r}), \quad r \in \mathbb{N}.$$

The upshot

Since

$$\mathrm{Kar}(\Phi): \mathrm{Kar}(\mathcal{F}_{7/3,26}) \rightarrow \mathfrak{f}\text{-mod}$$

is full and essentially surjective,

$\mathfrak{f}\text{-mod}$ is a quotient of $\mathrm{Kar}(\mathcal{F}_{7/3,26})$.

Philosophy

All homomorphisms between f.d. \mathfrak{f} -modules are built from the “flip” symmetry, the bilinear form, and a trilinear form (or the multiplication in the Albert algebra).

Question

Is the functor $\mathrm{Kar}(\Phi)$ faithful, and hence an equivalence of categories, or are there further relations?

Summary

- \mathcal{F} is a type F_4 analogue of the unoriented and oriented Brauer categories.
- $\text{End}_{\mathcal{F}}(\mathbb{1}^{\otimes r})$ is a type F_4 analogue of the oriented (walled) and unoriented Brauer algebras.
- We now have diagrammatic tools to study representation theory in type F_4 .

Further directions

Quantum analogue

There are **quantum** diagrammatic categories:

- **Type A** : framed HOMFLYPT skein category
- **Types BCD** : Kauffman skein category
- **Type G_2** : Kuperberg (1994)

There should exist a quantum analogue of the category \mathcal{F} .

Webs

- Webs given a diagrammatic description of the category generated by the fundamental modules.
- First described in **rank 2** by Kuperberg (1994, 1996).
- Arbitrary **type A** by Cautis–Kamnizter–Morrison (2014) and type C by Bodish–Elias–Rose–Tatham (2021).
- \mathcal{F} is a starting point for developing **F_4 webs**.