

Categorical comultiplication

$$\Delta(\times) = \begin{array}{c} \times \\ \text{blue} \end{array} + \begin{array}{c} \times \\ \text{red} \end{array} + \begin{array}{c} \times \\ \text{blue} \end{array} + \begin{array}{c} \times \\ \text{red} \end{array}.$$

Alistair Savage
University of Ottawa

Slides available online: alistairsavage.ca/talks

Based on joint work with Jon Brundan and Ben Webster.

Outline

Goal: Discuss concept of categorical comultiplication, with examples.

Overview:

- 1 Motivation: Why categorical comultiplication?
- 2 Example 1: The symmetric group category
- 3 Example 2: The affine Hecke algebra category
- 4 Example 3: The quantum Heisenberg category

Motivation

Suppose R is an associative \mathbb{C} -algebra and M, N are R -modules.

We have the **outer tensor product** $R \otimes R$ -module $M \boxtimes N$:

$$(r \otimes s)(m \otimes n) = rm \otimes sn.$$

If we want an **internal tensor product** R -module $M \otimes N$, we need a **comultiplication**

$$\delta: R \rightarrow R \otimes R,$$

so that we can define the R -action

$$r \cdot (m \otimes n) = \delta(r)(m \otimes n).$$

Motivating observation:

- If \mathfrak{g} is a f.d. simple Lie algebra, then $x \in U(\mathfrak{g})$ annihilates every f.d. simple module iff $x = 0$.
- For arbitrary Kac–Moody algebras this is false. But it's true if we consider (highest weight) \otimes (lowest weight) modules.

Strict monoidal categories

A **strict monoidal category** is a category \mathcal{C} equipped with

- a bifunctor (the **tensor product**) $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and
- a **unit object** $\mathbb{1}$,

such that, for objects A, B, C and morphisms f, g, h ,

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$,
- $(f \otimes g) \otimes h = f \otimes (g \otimes h)$,
- $1_{\mathbb{1}} \otimes f = f = f \otimes 1_{\mathbb{1}}$.

A **strict \mathbb{C} -linear monoidal category** is a strict monoidal category such that

- each morphism space is a \mathbb{C} -module,
- composition of morphisms is \mathbb{C} -bilinear,
- tensor product of morphisms is \mathbb{C} -bilinear.

Motivation

Suppose \mathcal{C} is a strict \mathbb{C} -linear monoidal category.

A **module category** over \mathcal{C} is a \mathbb{C} -linear category \mathcal{V} , together with a \mathbb{C} -linear monoidal functor

$$\mathcal{C} \rightarrow \mathcal{E}nd(\mathcal{V}),$$

where $\mathcal{E}nd(\mathcal{V})$ is the strict \mathbb{C} -linear monoidal category of \mathbb{C} -linear endofunctors and natural transformations.

If \mathcal{V}, \mathcal{W} are \mathcal{C} -module categories, then $\mathcal{V} \boxtimes \mathcal{W}$ is a module category over the **symmetric product category** $\mathcal{C} \odot \mathcal{C}$.

To form the tensor product of \mathcal{C} -module categories, we need a functor

$$\Delta: \mathcal{C} \rightarrow \mathcal{C} \odot \mathcal{C}.$$

Goal

See some interesting examples of such **categorical comultiplications**.

String diagrams for strict monoidal categories

We use **string diagrams** to represent morphisms:

$$f: A \rightarrow B \rightsquigarrow \begin{array}{c} B \\ | \\ \textcircled{f} \\ | \\ A \end{array}, \quad 1_A: A \rightarrow A \rightsquigarrow \begin{array}{c} A \\ | \\ \text{---} \\ | \\ A \end{array}$$

Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:

$$\begin{array}{c} \textcircled{f} \\ | \\ \textcircled{g} \end{array} = \textcircled{fg} \quad \textcircled{f} \otimes \textcircled{g} = \begin{array}{c} \textcircled{f} \quad \textcircled{g} \end{array}$$

We have the **interchange law**:

$$\begin{array}{c} \textcircled{f} \\ | \\ \text{---} \\ | \\ \textcircled{g} \end{array} = \begin{array}{c} \textcircled{f} \quad \textcircled{g} \end{array} = \begin{array}{c} \text{---} \\ | \\ \textcircled{f} \quad \textcircled{g} \end{array}$$

The symmetric group category

Let Sym be the free \mathbb{C} -linear symmetric monoidal category on one object. So Sym has one generating object I and we denote

$$1_I = |$$

We have one generating morphism

$$\times : I \otimes I \rightarrow I \otimes I.$$

We impose the relations:

$$\begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} = \begin{array}{c} | \\ | \end{array}, \quad \begin{array}{c} \diagup \diagdown \\ \diagup \diagdown \end{array} = \begin{array}{c} \diagdown \diagup \\ \diagdown \diagup \end{array}.$$

We have that

$$\text{End}_{Sym}(I^{\otimes n}) = \mathbb{k}S_n$$

is the group algebra of the **symmetric group** on n letters.

Categorification of symmetric functions

We have

$$K_0(\text{Kar}(\mathcal{S}ym)) \cong \bigoplus_{n=0}^{\infty} K_0(\mathbb{k}S_n) \cong \text{Sym}, \quad [S_\lambda] \mapsto s_\lambda,$$

where, for a partition λ ,

$$S_\lambda = \text{Specht module}, \quad s_\lambda = \text{Schur function}.$$

Multiplication on $\mathcal{S}ym$ -modules is “induction product”

$$X \circ Y = \text{Ind}_{S_m \times S_n}^{S_{m+n}} X \boxtimes Y.$$

This is a **categorification** of the ring of symmetric functions.

Question: Can we categorify the well-known coproduct on symmetric functions?

Symmetric product of monoidal categories

The **symmetric product** of strict \mathbb{C} -linear monoidal categories $\mathcal{C} \odot \mathcal{D}$ is obtained by:

- 1 taking the free product of \mathcal{C} and \mathcal{D} , i.e. the strict \mathbb{C} -linear monoidal category defined by the disjoint union of the given generators and relations of \mathcal{C} and \mathcal{D} ;
- 2 adjoining isomorphisms $\sigma_{X,Y}: X \otimes Y \xrightarrow{\cong} Y \otimes X$ for each $X \in \mathcal{C}$, $Y \in \mathcal{D}$, subject to the relations

$$\begin{aligned}\sigma_{X_1 \otimes X_2, Y} &= (\sigma_{X_1, Y} \otimes 1_{X_2}) \circ (1_{X_1} \otimes \sigma_{X_2, Y}), \\ \sigma_{X, Y_1 \otimes Y_2} &= (1_{Y_1} \otimes \sigma_{X, Y_2}) \circ (\sigma_{X, Y_1} \otimes 1_{Y_2}), \\ \sigma_{X_2, Y} \circ (f \otimes 1_Y) &= (1_Y \otimes f) \circ \sigma_{X_1, Y}, \\ \sigma_{X, Y_2} \circ (1_X \otimes g) &= (g \otimes 1_X) \circ \sigma_{X, Y_1},\end{aligned}$$

for all $X, X_1, X_2 \in \mathcal{C}$, $Y, Y_1, Y_2 \in \mathcal{D}$ and $f: X_1 \rightarrow X_2$, $g: Y_1 \rightarrow Y_2$.

We have

$$K_0(\mathcal{C} \odot \mathcal{D}) \cong K_0(\mathcal{C}) \otimes K_0(\mathcal{D}).$$

$Sym \odot Sym$

The symmetric product $Sym \odot Sym$ has two generating objects,

$$| \quad \text{and} \quad |,$$

and generating morphisms

$$\times, \quad \times, \quad \sigma_{|,|} = \times, \quad \sigma_{|,|}^{-1} = \times,$$

subject to the relations

$$\begin{array}{cccc} \text{blue crossing} = ||, & \text{blue crossing} = \text{blue crossing}, & \text{red crossing} = ||, & \text{red crossing} = \text{red crossing}, \end{array}$$

$$\begin{array}{cccc} \text{blue/red crossing} = \text{blue/red crossing}, & \text{red/blue crossing} = \text{red/blue crossing}, & \text{blue crossing} = ||, & \text{red crossing} = ||. \end{array}$$

Categorical comultiplication on Sym

Consider the strict \mathbb{C} -linear monoidal functor

$$\begin{aligned} \Delta: Sym &\rightarrow \text{Add}(Sym \odot Sym), \\ I &\mapsto I \oplus I, \quad \times \mapsto \color{blue}{\times} + \color{red}{\times} + \color{blue}{\times} + \color{red}{\times}. \end{aligned}$$

This induces a strict \mathbb{C} -linear monoidal functor

$$\Delta: \text{Kar}(Sym) \rightarrow \text{Kar}(Sym \odot Sym).$$

Passing to split Grothendieck groups yields

$$\begin{array}{ccc} K_0(Sym) & \xrightarrow{[\Delta]} & K_0(Sym \odot Sym) \\ \downarrow \cong & & \downarrow \cong \\ Sym & \xrightarrow{\delta} & Sym \otimes Sym \end{array}$$

where δ is the usual comultiplication on Sym :

$$\delta(h_n) = \sum_{r=0}^n h_{n-r} \otimes h_r, \quad \delta(e_n) = \sum_{r=0}^n e_{n-r} \otimes e_r, \quad \delta(p_n) = p_n \otimes 1 + 1 \otimes p_n.$$

Hecke algebras

Fix $q \in \mathbb{C}^\times$. Let $\mathcal{H}(q)$ be the strict \mathbb{C} -linear monoidal category with one generating object \uparrow , generating morphisms

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array}, \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow$$

and relations

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array} = \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array} \begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array}, \\ \begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} - \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array} = (q - q^{-1}) \begin{array}{c} \uparrow \\ \uparrow \end{array} \quad (\text{Conway skein relation}).$$

Then

$$H_n(q) := \text{End}_{\mathcal{H}(q)}(\uparrow^{\otimes n})$$

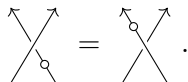
is the **Iwahori–Hecke algebra** of type A_{n-1} .

Affine Hecke algebras

Define $\mathcal{H}^{\text{aff}}(q)$ by adjoining to $\mathcal{H}(q)$ the **invertible** morphism

$$\uparrow\!:\! \uparrow \rightarrow \uparrow$$

and relation



The diagram shows an equality between two crossings. On the left, a crossing of two lines has a small circle (dot) on the lower-left strand. On the right, a crossing of two lines has a small circle (dot) on the upper-left strand. The two diagrams are separated by an equals sign.

Then

$$H_n^{\text{aff}}(q) := \text{End}_{\mathcal{H}^{\text{aff}}(q)}(\uparrow^{\otimes n})$$

is the **affine Hecke algebra** of type A_{n-1} .

Comultiplication on $\mathcal{H}^{\text{aff}}(q)$

We'd like to define a categorical comultiplication:

$$\Delta: \mathcal{H}^{\text{aff}}(q) \rightarrow \text{Add}(\mathcal{H}^{\text{aff}}(q) \odot \mathcal{H}^{\text{aff}}(q)).$$

However, naive guesses such as

$$\begin{aligned} \uparrow &\mapsto \uparrow \oplus \uparrow, & \uparrow \circlearrowleft &\mapsto \uparrow \circlearrowleft + \uparrow \circlearrowleft, \\ \begin{array}{c} \nearrow \\ \nwarrow \end{array} &\mapsto \begin{array}{c} \nearrow \\ \nwarrow \end{array} + \begin{array}{c} \nwarrow \\ \nearrow \end{array} + \begin{array}{c} \nwarrow \\ \nearrow \end{array} + \begin{array}{c} \nearrow \\ \nwarrow \end{array}, & \begin{array}{c} \nearrow \\ \nwarrow \end{array} &\mapsto \begin{array}{c} \nearrow \\ \nwarrow \end{array} + \begin{array}{c} \nwarrow \\ \nearrow \end{array} + \begin{array}{c} \nwarrow \\ \nearrow \end{array} + \begin{array}{c} \nearrow \\ \nwarrow \end{array}, \end{aligned}$$

do **not** work! They don't respect the relations.

It turns out we'll have to **localize** to define the categorical comultiplication.

The neutral crossing

Consider the **neutral crossing**

$$\overline{\times} := \times - {}_{-1}\overline{\times} = \times - \overline{\times}^{-1} = \times - \overline{\times}^{-1} + \uparrow\uparrow$$

(Related to **intertwining operators** for affine Hecke algebras.) Direct computation shows that

$$\begin{aligned} \overline{\times} &= \overline{\times}^{\circlearrowleft}, & \overline{\times} &= \overline{\times}^{\circlearrowright}, & \overline{\times} &= \overline{\times}, \\ \overline{\times} &= \left(q \uparrow\uparrow - q^{-1} {}_{-1}\overline{\times} \right) \circ \left(q \uparrow\uparrow - q^{-1} \overline{\times} \right). \end{aligned}$$

We'd like to work with a new presentation of $\mathcal{H}^{\text{aff}}(q)$ involving $\overline{\times}$.

To do this, we need to be able to express $\overline{\times}$ and $\overline{\times}$ in terms of \times . This means we need to invert

$$\uparrow\uparrow - \overline{\times}^{-1}$$

A new presentation

Define $\overline{\mathcal{H}^{\text{aff}}}(q)$ by adjoining to $\mathcal{H}^{\text{aff}}(q)$ a morphism

$$\begin{array}{c} \uparrow \leftarrow \uparrow \\ \uparrow \quad \uparrow \end{array} := \left(\begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} - \begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} \right)^{-1}. \quad (\boxtimes)$$

Proposition

The strict \mathbb{C} -linear monoidal category $\overline{\mathcal{H}^{\text{aff}}}(q)$ is generated by one object \uparrow and morphisms $\begin{array}{c} \nearrow \\ \circ \\ \searrow \end{array}$, $\begin{array}{c} \uparrow \\ \circ \end{array}$, $\begin{array}{c} \uparrow \leftarrow \uparrow \\ \uparrow \quad \uparrow \end{array}$, subject to the relation (\boxtimes) and

$$\begin{array}{c} \nearrow \\ \circ \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \searrow \\ \circ \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \circ \end{array} = \begin{array}{c} \circ \\ \nearrow \\ \searrow \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \searrow \\ \searrow \\ \nearrow \end{array},$$

$$\begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \left(q \begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} \right) \left(q \begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} \right)^{-1}.$$

Sketch of proof: $\begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \uparrow \leftarrow \uparrow \\ \uparrow \quad \uparrow \end{array} - \begin{array}{c} \uparrow \leftarrow \uparrow \\ \uparrow \quad \uparrow \end{array}$

Comultiplication on $\mathcal{H}^{\text{aff}}(q)$

Recall:

$$\begin{array}{c} \diagup \circ \diagdown \\ \diagdown \circ \diagup \end{array} = \begin{array}{c} \diagup \\ \diagdown \end{array} = \left(q \begin{array}{c} \uparrow \\ \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \\ \circ \\ \uparrow \end{array} \right) \left(q \begin{array}{c} \uparrow \\ \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \\ \circ \\ \uparrow \end{array} - 1 \right).$$

However the two-colored crossings behave differently:

$$\begin{array}{c} \diagup \\ \diagdown \end{array} \circ \begin{array}{c} \diagdown \\ \diagup \end{array} = \begin{array}{c} \diagup \\ \diagdown \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array} \quad \text{and} \quad \begin{array}{c} \diagdown \\ \diagup \end{array} \circ \begin{array}{c} \diagup \\ \diagdown \end{array} = \begin{array}{c} \diagdown \\ \diagup \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}.$$

We can “fix” this:

$$\begin{aligned} & \left(q \begin{array}{c} \diagup \\ \diagdown \end{array} - q^{-1} \begin{array}{c} \diagup \\ \circ \\ \diagdown \end{array} \right) \circ \left(q \begin{array}{c} \diagdown \\ \diagup \end{array} - q^{-1} \begin{array}{c} \diagdown \\ \circ \\ \diagup \end{array} - 1 \right) \\ &= \left(q \begin{array}{c} \uparrow \\ \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \\ \circ \\ \uparrow \end{array} \right) \left(q \begin{array}{c} \uparrow \\ \uparrow \end{array} - q^{-1} \begin{array}{c} \uparrow \\ \circ \\ \uparrow \end{array} - 1 \right). \end{aligned}$$

Comultiplication on $\mathcal{H}^{\text{aff}}(q)$

Let $\mathcal{H}^{\text{aff}}(q) \overline{\odot} \mathcal{H}^{\text{aff}}(q)$ denote the symmetric product $\mathcal{H}^{\text{aff}}(q) \odot \mathcal{H}^{\text{aff}}(q)$ where we have adjoined an inverse

$$\begin{array}{c} \uparrow \leftarrow \uparrow \\ \circ \quad \circ \end{array} := \left(\begin{array}{c} \uparrow \quad \uparrow \\ \circ \quad \circ \end{array} - \begin{array}{c} \uparrow \quad \uparrow \\ \circ \quad \circ \end{array} -1 \right)^{-1}.$$

Proposition

There is a strict \mathbb{C} -linear monoidal functor

$$\Delta: \overline{\mathcal{H}^{\text{aff}}(q)} \rightarrow \text{Add} \left(\mathcal{H}^{\text{aff}}(q) \overline{\odot} \mathcal{H}^{\text{aff}}(q) \right)$$

given on objects by $\uparrow \mapsto \uparrow \oplus \uparrow$, and on morphisms by

$$\begin{array}{c} \nearrow \searrow \\ \circ \quad \circ \end{array} \mapsto \begin{array}{c} \nearrow \searrow \\ \circ \quad \circ \end{array} + \begin{array}{c} \searrow \nearrow \\ \circ \quad \circ \end{array} + \left(q \begin{array}{c} \nearrow \searrow \\ \circ \quad \circ \end{array} - q^{-1} \begin{array}{c} \searrow \nearrow \\ \circ \quad \circ \end{array} -1 \right) + \left(q \begin{array}{c} \searrow \nearrow \\ \circ \quad \circ \end{array} - q^{-1} \begin{array}{c} \nearrow \searrow \\ \circ \quad \circ \end{array} -1 \right).$$

Quantum Heisenberg category

Fix $k \in \mathbb{Z}$. The **quantum Heisenberg category** (Licata–S., Brundan–S.–Webster) is defined by adjoining to $\mathcal{H}^{\text{aff}}(q)$ an object

↓

and morphisms and relations so that

- ↓ is right dual to ↑, meaning we have morphisms \cup and \cap such that

$$\cup = \uparrow, \quad \cap = \downarrow;$$

- we have an isomorphism

$$\begin{aligned} \uparrow \otimes \downarrow &\cong \downarrow \otimes \uparrow \oplus \mathbb{1}^{\oplus k} \quad (\text{when } k \geq 0), \\ \uparrow \otimes \downarrow \oplus \mathbb{1}^{\oplus (-k)} &\cong \downarrow \otimes \uparrow \quad (\text{when } k \leq 0) \end{aligned}$$

(the **inversion relation**).

Acts on modules for cyclotomic Hecke algebras.

Comultiplication on the quantum Heisenberg category

We can extend the categorical comultiplication on $\mathcal{H}^{\text{aff}}(q)$ to one on the quantum Heisenberg category, with

$$\downarrow \mapsto \downarrow \oplus \downarrow, \quad \cap \mapsto \cap + \cap, \quad \cup \mapsto \cup + \cup.$$

This allows us to take tensor products of module categories.

In particular, we can form **generalized cyclotomic quotients**, which are large representations on which the quantum Heisenberg category acts faithfully.

We use these large representations to prove a **basis theorem** for the quantum Heisenberg category.

Technique can be used for other categories: degenerate Heisenberg category, categorified quantized enveloping algebras,...