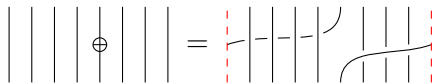


Affinization of monoidal categories

(Monoidal categories on cylinders)



Alistair Savage
University of Ottawa

Slides: alistsavsage.ca/talks

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Outline

Goal: Guided by intuition of monoidal categories on cylinders, formalize the notion of **affinization** of a monoidal category.

Overview:

- 1 String diagrams for strict monoidal categories
- 2 Examples
- 3 Definition of affinization
- 4 Dot presentation
- 5 Examples
- 6 Horizontal and vertical trace

Strict monoidal categories

A **strict monoidal category** is a category \mathcal{C} equipped with

- a bifunctor (the **tensor product**) $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and
- a **unit object** $\mathbb{1}$,

such that, for objects A, B, C and morphisms f, g, h ,

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$,
- $(f \otimes g) \otimes h = f \otimes (g \otimes h)$,
- $1_{\mathbb{1}} \otimes f = f = f \otimes 1_{\mathbb{1}}$.

Remark: Non-strict monoidal categories

In a (not necessarily strict) **monoidal category**, the equalities above are replaced by isomorphism, and we impose some **coherence conditions**.

Every monoidal category is monoidally equivalent to a strict one.

\mathbb{k} -linear monoidal categories

Fix a commutative ground ring \mathbb{k} .

A **strict \mathbb{k} -linear monoidal category** is a strict monoidal category such that

- each morphism space is a \mathbb{k} -module,
- composition of morphisms is \mathbb{k} -bilinear,
- tensor product of morphisms is \mathbb{k} -bilinear.

The interchange law

The axioms of a strict monoidal category imply the **interchange law**: For $A_1 \xrightarrow{f} A_2$ and $B_1 \xrightarrow{g} B_2$, the following diagram commutes:

$$\begin{array}{ccc} A_1 \otimes B_1 & \xrightarrow{1 \otimes g} & A_1 \otimes B_2 \\ f \otimes 1 \downarrow & \searrow f \otimes g & \downarrow f \otimes 1 \\ A_2 \otimes B_1 & \xrightarrow{1 \otimes g} & A_2 \otimes B_2 \end{array}$$

String diagrams

Fix a strict monoidal category \mathcal{C} .

We will denote a morphism $f: A \rightarrow B$ by:



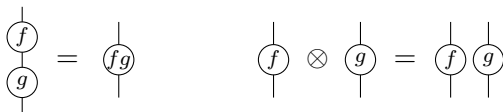
The **identity map** $1_A: A \rightarrow A$ is a string with no label:



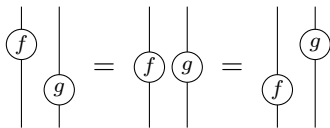
We sometimes omit the object labels when they are clear or unimportant.

String diagrams

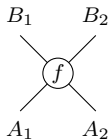
Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:



The **interchange law** then becomes:



A morphism $f: A_1 \otimes A_2 \rightarrow B_1 \otimes B_2$ can be depicted:



Example: monoidally generated symmetric groups

Define a strict monoidal category Sym with one generating object \uparrow and denote

$$1_{\uparrow} = \uparrow$$

We have one generating morphism

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow .$$

We impose the relations:

$$\begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \searrow \\ \searrow \\ \nearrow \end{array} .$$

Then

$$\text{End}_{Sym}(\uparrow^{\otimes n}) = S_n$$

is the **symmetric group** on n letters.

Example: monoidally generated symmetric groups

This monoidal presentation of S_n is very efficient! We only needed

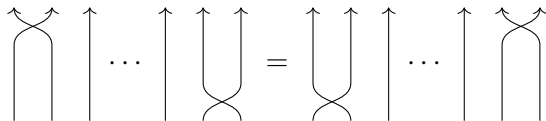
- one generating morphism, and
- two relations,

to get **all** the symmetric groups.

Note that the “distant braid relation”

$$s_i s_j = s_j s_i, \quad |i - j| > 1$$

for simple transpositions follows for free from the interchange law:



Note: If we define Sym to be \mathbb{k} -linear, then $\text{End}_{Sym}(\uparrow^{\otimes n}) = \mathbb{k}S_n$.

Braidings

A strict monoidal category \mathcal{C} is **braided** if it is equipped with isomorphisms

$$\beta_{X,Y} = \begin{array}{c} \diagup \quad \diagdown \\ X \quad Y \end{array} : X \otimes Y \rightarrow Y \otimes X, \quad \beta_{X,Y}^{-1} = \begin{array}{c} \diagdown \quad \diagup \\ Y \quad X \end{array},$$

for all $X, Y \in \text{Ob}(\mathcal{C})$, satisfying

$$\begin{array}{c} \diagup \quad \diagdown \\ X \quad Y \otimes Z \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ X \quad Y \quad Z \end{array}, \quad \begin{array}{c} \diagdown \quad \diagup \\ X \otimes Y \quad Z \end{array} = \begin{array}{c} \diagup \quad \diagdown \\ X \quad Y \quad Z \end{array},$$

and

$$\begin{array}{c} \diagup \quad \diagdown \\ \text{---} \quad \text{---} \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \text{---} \quad \text{---} \end{array}, \quad \begin{array}{c} \diagdown \quad \diagup \\ \text{---} \quad \text{---} \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \text{---} \quad \text{---} \end{array} \circ f, \quad \begin{array}{c} \diagdown \quad \diagup \\ \text{---} \quad \text{---} \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \text{---} \quad \text{---} \end{array} \circ f,$$

for all morphisms f in \mathcal{C} (and any compatible labelling of strands).

Braid groups

The category *Braid* of braids over the disc is isomorphic to the strict monoidal category generated by a single object \uparrow , and morphisms

$$\begin{array}{c} \nearrow \\ \nwarrow \end{array}, \begin{array}{c} \nwarrow \\ \nearrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow,$$

subject to the relations

$$\begin{array}{c} \nearrow \\ \nwarrow \end{array} \begin{array}{c} \nwarrow \\ \nearrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array} = \begin{array}{c} \nwarrow \\ \nearrow \end{array} \begin{array}{c} \nearrow \\ \nwarrow \end{array}, \quad \begin{array}{c} \nwarrow \\ \nearrow \end{array} \begin{array}{c} \nearrow \\ \nwarrow \end{array} = \begin{array}{c} \nearrow \\ \nwarrow \end{array} \begin{array}{c} \nwarrow \\ \nearrow \end{array}.$$

Universal property: *Braid* is the free braided monoidal category generated by a single object.

Note that

$$\text{End}_{\text{Braid}}(\uparrow^{\otimes n}) = \text{braid group of type } A_{n-1}.$$

Hecke algebras

Let $\mathit{Braid}_{\mathbb{k}}$ be the \mathbb{k} -linearization of Braid , i.e. the strict \mathbb{k} -linear monoidal category with the same generators and relations.

Fix $z \in \mathbb{k}^{\times}$, and let \mathcal{H} be obtained from $\mathit{Braid}_{\mathbb{k}}$ by imposing the **Conway skein relation**

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} - \begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array}.$$

Then

$$\mathrm{End}_{\mathcal{H}}(\uparrow^{\otimes n}) = \text{Iwahori-Hecke algebra of type } A_{n-1}.$$

(Often one sets $z = q - q^{-1}$ for some $q \in \mathbb{k}^{\times}$.)

Duals

Suppose a strict monoidal category \mathcal{C} has two objects \uparrow and \downarrow , with

$$1_{\uparrow} = \uparrow \quad , \quad 1_{\downarrow} = \downarrow .$$

A morphism $\mathbb{1} \rightarrow \downarrow \otimes \uparrow$ would have string diagram

$$\begin{array}{c} \swarrow \quad \searrow \\ \downarrow \text{ (dotted)} \end{array} \quad , \quad \text{where} \quad \begin{array}{c} \vdots \\ \vdots \\ \vdots \end{array} = 1_{\mathbb{1}} .$$

We typically omit the dotted line and draw:

$$\cup : \mathbb{1} \rightarrow \downarrow \otimes \uparrow .$$

Similarly, we can have

$$\cap : \uparrow \otimes \downarrow \rightarrow \mathbb{1} .$$

Duals

We say that \downarrow is **right dual** to \uparrow (and \uparrow is **left dual** to \downarrow) if there exist morphisms

$$\cup : \mathbb{1} \rightarrow \downarrow \otimes \uparrow \quad \text{and} \quad \cap : \uparrow \otimes \downarrow \rightarrow \mathbb{1}.$$

such that

$$\begin{array}{c} \downarrow \\ \cup \\ \downarrow \end{array} = \downarrow \quad \text{and} \quad \begin{array}{c} \uparrow \\ \cap \\ \uparrow \end{array} = \uparrow.$$

Objects \uparrow and \downarrow are both left and right dual to each other if we also have

$$\cup : \mathbb{1} \rightarrow \uparrow \otimes \downarrow \quad \text{and} \quad \cap : \downarrow \otimes \uparrow \rightarrow \mathbb{1}$$

such that

$$\begin{array}{c} \uparrow \\ \cup \\ \uparrow \end{array} = \uparrow \quad \text{and} \quad \begin{array}{c} \downarrow \\ \cap \\ \downarrow \end{array} = \downarrow.$$

We say \mathcal{C} is **rigid** if all objects have left and right duals.

Duals: example

Let \mathbb{k} be a field and consider the category $\text{Vect}_{\mathbb{k}}$ of f.d. \mathbb{k} -vector spaces.

Unit object: \mathbb{k}

Fix a f.d. \mathbb{k} -vector space V .

Claim: The dual vector space V^* is both left and right dual to V , in the sense mentioned above.

Proof: Fix a basis B of V . Let $\{\delta_v : v \in B\}$ be the dual basis of V^* . Viewing V and \uparrow and V^* as \downarrow , we define

$$\cup : \mathbb{k} \rightarrow V^* \otimes V,$$

$$\cap : V \otimes V^* \rightarrow \mathbb{k},$$

$$\cup : \mathbb{k} \rightarrow V \otimes V^*,$$

$$\cap : V^* \otimes V \rightarrow \mathbb{k},$$

$$1 \mapsto \sum_{v \in B} \delta_v \otimes v,$$

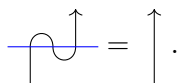
$$v \otimes f \mapsto f(v),$$

$$1 \mapsto \sum_{v \in B} v \otimes \delta_v,$$

$$f \otimes v \mapsto f(v).$$

Duals: example

Let's check the relation


$$\text{wavy line} = \text{vertical arrow}.$$

The left-hand side is the composition

$$V \cong V \otimes \mathbb{k} \xrightarrow{1_V \otimes \cup} V \otimes V^* \otimes V \xrightarrow{\cap \otimes 1_V} \mathbb{k} \otimes V \cong V,$$
$$w \mapsto w \otimes 1 \mapsto \sum_{v \in B} w \otimes \delta_v \otimes v \mapsto \sum_{v \in B} \delta_v(w) \otimes v \mapsto \sum_{v \in V} \delta_v(w) v = w.$$

The verification of the other relations is analogous.

Self-dual objects

An object I is **self-dual** if there exist morphisms

$$\cup : \mathbb{1} \rightarrow I \otimes I. \quad \text{and} \quad \cap : I \otimes I \rightarrow \mathbb{1}.$$

such that

$$\cup \cap = \text{id} = \cap \cup.$$

Example

In the category of **finite-dimensional inner product spaces** (over \mathbb{R} or \mathbb{C}), all objects V are self-dual.

$$\begin{aligned} \cap : V \otimes V &\rightarrow \mathbb{k}, & v \otimes w &\mapsto \langle v, w \rangle, \\ \cup : \mathbb{k} &\rightarrow V \otimes V, & 1 &\mapsto \sum_{v \in B} v \otimes v, \end{aligned}$$

where B is an orthonormal basis for V .

Pivotal categories

A strict monoidal category \mathcal{C} is a **strict pivotal category** if every object X has a right dual X^\vee and the following conditions are satisfied:

- 1 For all objects X and Y in \mathcal{C} ,

$$(X^\vee)^\vee = X, \quad (X \otimes Y)^\vee = Y^\vee \otimes X^\vee, \quad \mathbf{1}^\vee = \mathbf{1},$$

$X \otimes Y = XY$ and $X \otimes Y = XY$.

- 2 For every morphism $f: X \rightarrow Y$ in \mathcal{C} ,

$X \xrightarrow{f} Y = X \xrightarrow{f} Y : Y^\vee \rightarrow X^\vee$.

Intuition: In a strict pivotal category, morphisms are invariant under isotopy.

Framed tangles

The category \mathcal{FT} of **framed tangles** over the disc is generated by a single object I and morphisms

$$\times, \times, \cup, \cap,$$

subject to the relations

$$\begin{aligned} \cup = \cap &= \cup, & \cap = \cap, & \cup = \cup, \\ d_1 = d_2, & \cup = \cup = \cup, & \cup = \cup. & \end{aligned}$$

Universal property: \mathcal{FT} is the free ribbon category generated by a self-dual object.

Kauffman skein category

Fix $z, t \in \mathbb{k}^\times$. The **Kauffman skein category** $\mathcal{KS}(z, t)$ is the category obtained from $\mathcal{FT}_{\mathbb{k}}$ by imposing the relations

$$\begin{array}{c} \diagdown \\ \diagup \end{array} + \begin{array}{c} \diagup \\ \diagdown \end{array} = z \left| \begin{array}{c} | \\ | \end{array} \right. + z \begin{array}{c} \cup \\ \cap \end{array} \quad (\text{Kauffman skein relation})$$

and

$$\begin{array}{c} \circlearrowleft \end{array} = t \left| \begin{array}{c} | \\ | \end{array} \right., \quad \bigcirc = \left(1 - \frac{t + t^{-1}}{z} \right) \mathbb{1}_{\mathbb{1}}.$$

Then

$$\begin{aligned} \text{End}_{\mathcal{KS}(z,t)}(\mathbb{1}^{\otimes n}) &\cong \text{Kauffman tangle algebra} \\ &\cong \text{Birman–Murakami–Wenzl (BMW) algebra.} \end{aligned}$$

Viewing a link as an element of $\text{End}_{\mathcal{KS}(z,t)}(\mathbb{1})$ gives rise to its **Kauffman polynomial**.

Affinization

Definition

The **affinization** of a strict monoidal category \mathcal{C} is the category $\text{Aff}(\mathcal{C})$ obtained from \mathcal{C} by adjoining invertible morphisms

$$\xi_{X,Y}: X \otimes Y \rightarrow Y \otimes X, \quad X, Y \in \text{Ob}(\mathcal{C}),$$

subject to the relations

$$\xi_{X,Y \otimes Z} = \xi_{Z \otimes X, Y} \circ \xi_{X \otimes Y, Z}, \quad \xi_{X_2, Y_2} \circ (g \otimes f) = (f \otimes g) \circ \xi_{X_1, Y_1}.$$

for all $X, Y, Z \in \text{Ob}(\mathcal{C})$, $f \in \text{Hom}_{\mathcal{C}}(Y_1, Y_2)$, and $g \in \text{Hom}_{\mathcal{C}}(X_1, X_2)$.

Picture morphisms in $\text{Aff}(\mathcal{C})$ as \mathcal{C} -diagrams on the cylinder, with

$$\xi_{X,Y} = \begin{array}{c} \text{Cylinder diagram with two strands: } X \text{ (left) and } Y \text{ (right). The strands cross twice, with } X \text{ on top and } Y \text{ on bottom in the middle section.} \\ X \quad Y \end{array}, \quad \xi_{X,Y}^{-1} = \begin{array}{c} \text{Cylinder diagram with two strands: } Y \text{ (left) and } X \text{ (right). The strands cross twice, with } Y \text{ on top and } X \text{ on bottom in the middle section.} \\ Y \quad X \end{array}. \quad \text{"coils"}$$

Affinization

Cutting open the cylinder, we draw

$$\xi_{X,Y} = \begin{array}{c} \text{Cylinder with } X \text{ and } Y \text{ strands} \\ \hline \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ X \quad Y \end{array} \end{array}, \quad \xi_{X,Y}^{-1} = \begin{array}{c} \text{Cylinder with } X \text{ and } Y \text{ strands} \\ \hline \begin{array}{c} X \quad Y \\ \diagup \quad \diagdown \\ Y \quad X \end{array} \end{array}.$$

Then the defining relations become

$$\begin{array}{c} Y \otimes Z \quad X \\ \diagdown \quad / \\ X \quad Y \otimes Z \end{array} = \begin{array}{c} YZ \quad X \\ \diagdown \quad / \\ X \quad YZ \end{array}, \quad \begin{array}{c} Y_2 \quad X_2 \\ \diagdown \quad / \\ X_1 \quad Y_1 \end{array} = \begin{array}{c} Y_2 \quad X_2 \\ \oplus \quad \oplus \\ X_1 \quad Y_1 \end{array}.$$

If \mathcal{C} is **braided**, then $\text{Aff}(\mathcal{C})$ is **strict monoidal**, with tensor product given by nesting cylinders. For example,

$$\begin{array}{c} Y \quad X \\ \diagdown \quad / \\ X \quad Y \end{array} \otimes \begin{array}{c} | \\ Z \end{array} = \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ X \quad YZ \end{array}, \quad \begin{array}{c} | \\ X \end{array} \otimes \begin{array}{c} Z \quad Y \\ \diagdown \quad / \\ Y \quad Z \end{array} = \begin{array}{c} Z \quad Y \\ \diagdown \quad / \\ XY \quad Z \end{array}.$$

Dot presentation

Suppose \mathcal{C} is **braided**, so that $\text{Aff}(\mathcal{C})$ is strict monoidal. Define

$$\begin{array}{c} \oplus \\ X \end{array} := \xi_{\mathbb{1}, X} = \begin{array}{c} X \\ \text{[Cylinder with dot]} \\ X \end{array} = \begin{array}{c} X \\ \text{[Crossing]} \\ X \end{array}, \quad \begin{array}{c} \ominus \\ X \end{array} := \xi_{\mathbb{1}, X}^{-1} = \begin{array}{c} X \\ \text{[Cylinder with dot]} \\ X \end{array} = \begin{array}{c} X \\ \text{[Crossing]} \\ X \end{array}.$$

Thus $\begin{array}{c} | \\ | \\ | \\ | \\ \oplus \\ | \\ | \\ | \\ | \end{array} = \begin{array}{c} \text{[Crossing]} \\ | \\ | \\ | \\ | \end{array}, \quad \begin{array}{c} | \\ | \\ | \\ | \\ \ominus \\ | \\ | \\ | \\ | \end{array} = \begin{array}{c} \text{[Crossing]} \\ | \\ | \\ | \\ | \end{array}.$

Theorem

The **affinization** $\text{Aff}(\mathcal{C})$ is obtained from \mathcal{C} by adjoining morphisms

$$\begin{array}{c} \oplus \\ X \end{array}, \quad \begin{array}{c} \ominus \\ X \end{array} : X \rightarrow X, \quad X \in \text{Ob}(\mathcal{C}), \quad \text{subject to relations}$$

$$\begin{array}{c} \oplus \\ \oplus \end{array} = \begin{array}{c} \oplus \\ \oplus \end{array}, \quad \begin{array}{c} \oplus \\ X \otimes Y \end{array} = \begin{array}{c} \oplus \\ X \end{array} \begin{array}{c} \oplus \\ Y \end{array}, \quad \begin{array}{c} \oplus \\ f \end{array} = \begin{array}{c} f \\ \oplus \end{array}, \quad \begin{array}{c} \oplus \\ \oplus \end{array} = \begin{array}{c} \oplus \\ \oplus \end{array} = \begin{array}{c} | \end{array}.$$

Pivotal structures

Suppose \mathcal{C} is braided and **rigid/pivotal**. Then $\text{Aff}(\mathcal{C})$ is also rigid/pivotal.

Self-dual case (unoriented)

If X is self-dual, for strands labelled X we have

$$\oplus \cap = \cap \ominus, \quad \ominus \cap = \cap \oplus, \quad \oplus \cup = \cup \ominus, \quad \ominus \cup = \cup \oplus.$$

Non-self-dual case (oriented)

Define the **invertible dots**

$$\begin{array}{c} \uparrow \\ \circ \\ | \\ X \end{array} := \begin{array}{c} \uparrow \\ \oplus \\ | \\ X \end{array}, \quad \begin{array}{c} \downarrow \\ \circ \\ | \\ X \end{array} := \begin{array}{c} \downarrow \\ \ominus \\ | \\ X \end{array}.$$

Then we have

$$\begin{array}{c} \uparrow \\ \circ \\ | \\ \downarrow \end{array} = \begin{array}{c} \downarrow \\ \ominus \\ | \\ \uparrow \end{array}, \quad \begin{array}{c} \downarrow \\ \circ \\ | \\ \uparrow \end{array} = \begin{array}{c} \uparrow \\ \oplus \\ | \\ \downarrow \end{array}, \quad \begin{array}{c} \downarrow \\ \circ \\ | \\ \downarrow \end{array} = \begin{array}{c} \downarrow \\ \ominus \\ | \\ \downarrow \end{array}, \quad \begin{array}{c} \uparrow \\ \circ \\ | \\ \uparrow \end{array} = \begin{array}{c} \uparrow \\ \oplus \\ | \\ \uparrow \end{array}.$$

Examples: Towers of algebras

Suppose \mathcal{C} is a strict linear monoidal category such that

- objects are generated by a single object X ,
- $\text{Hom}_{\mathcal{C}}(X^{\otimes n}, X^{\otimes m}) = 0$ when $m \neq n$.

The collection $\mathcal{C}(n) := \text{End}_{\mathcal{C}}(X^{\otimes n})$, $n \in \mathbb{N}$, is a **tower of algebras**.

Braids

Category *Braid* of braids over the disc has

generators , , relations  =  = ,  = 

Affine braids

The **affinization** $\text{Aff}(\textit{Braid})$ is the category of braids over the annulus. Dot presentation obtained from *Braid* by adding

invertible generator , relation  = 

Examples: Towers of algebras

Hecke algebras

The category \mathcal{H} corresponding to the tower of **Hecke algebras of type A** is the \mathbb{k} -linearization of *Braid* modulo the Conway skein relation

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} - \begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} = (q - q^{-1}) \begin{array}{c} \uparrow \uparrow \end{array}, \quad q \in \mathbb{k}^\times.$$

Affine Hecke algebras

The **affinization** $\text{Aff}(\mathcal{H})$ corresponds to the tower of **affine Hecke algebras of type A** . Obtained from \mathcal{H} by adding

invertible generator $\hat{\circ}$, relation $\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} = \begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \circ \end{array}$

Examples: Tangles

The category \mathcal{FT} of **framed tangles** over the disc is generated by a single object I and morphisms

$$\times, \succ, \cup, \cap,$$

subject to the relations

$$\begin{aligned} \cup &= | = \cap, & \cap &= \cap', & \cup &= \cup', \\ d_1 &= \rho, & \cup &= || = \cup, & \times &= \times'. \end{aligned}$$

The **affinization** $\text{Aff}(\mathcal{FT})$ is the category of framed tangles over the **annulus**.

$\text{Aff}(\mathcal{FT})$ is obtained from \mathcal{FT} adding morphisms \oplus, \ominus and relations

$$\begin{aligned} \times_{\oplus} &= \times_{\ominus}, & \cap_{\oplus} &= \cap_{\ominus}, & \oplus &= \ominus = |. \end{aligned}$$

Similar results hold for non-framed tangles and oriented (framed and non-framed) tangles.

Examples: Skein categories

Kauffman skein category

The affinization of the **Kauffman skein category** over the disc is the Kauffman skein category over the annulus.

Dot presentation matches the **affine Kauffman skein category** of Gao–Rui–Song.

Temperley–Lieb category

The affinization of the **Temperley–Lieb category** is the **affine Temperley–Lieb category**.

Actions

Suppose \mathcal{C} , \mathcal{M} are braided strict pivotal categories and $F: \mathcal{C} \rightarrow \mathcal{M}$ is a monoidal functor.

Then \mathcal{C} acts on \mathcal{M} via the action

$$X \cdot M = F(X) \otimes M, \quad f \cdot g = F(f) \otimes g.$$

Theorem

The above can be extended to an action of $\text{Aff}(\mathcal{C})$ on \mathcal{M} with

$$\xi_X \cdot g = \begin{array}{c} N \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ (g) \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ M \\ F(X) \end{array}, \quad \xi_X^{-1} \cdot g = \begin{array}{c} N \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ (g) \\ | \\ \text{---} \\ | \\ \text{---} \\ | \\ M \\ F(X) \end{array}.$$

Note: More generally, have action whenever \mathcal{C} is a strict monoidal category and \mathcal{M} is a **balanced** strict monoidal category.

Horizontal trace

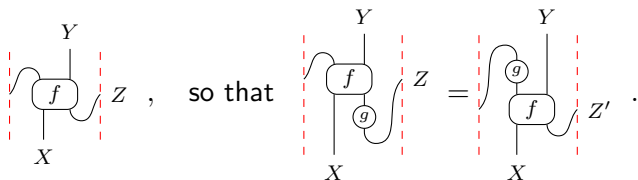
Suppose \mathcal{C} is a strict monoidal category. Fix $X, Y \in \text{Ob}(\mathcal{C})$ and consider pairs

$$(Z, f), \quad Z \in \text{Ob}(\mathcal{C}), \quad f: X \otimes Z \rightarrow Z \otimes Y.$$

Define an equivalence relation on such pairs by

$$(Z, f \circ (1_X \otimes g)) \sim (Z', (g \otimes 1_Y) \circ f), \\ g \in \text{Hom}_{\mathcal{C}}(Z, Z'), \quad f \in \text{Hom}_{\mathcal{C}}(X \otimes Z', Z \otimes Y).$$

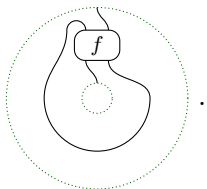
Picture the equivalence class $[Z, f]$ as the cylindrical diagram



Important: The cups and caps above have no precise meaning!

Horizontal trace

Also sometimes see $[Z, f]$ as the annular diagram



Important (again): The cups and caps above have no precise meaning!

Can define a natural composition, turning the horizontal trace $\text{Tr}_h(\mathcal{C})$ into a category.

If \mathcal{C} is **braided**, then the horizontal trace is a **monoidal category**.

Horizontal trace

Theorem

If \mathcal{C} is a **rigid** strict monoidal category then $\text{Aff}(\mathcal{C})$ is **isomorphic** to the horizontal trace $\text{Tr}_h(\mathcal{C})$ of \mathcal{C} .

In general (i.e. when \mathcal{C} is not rigid) the affinization and the horizontal trace are **different**.

It is the **affinization**, not the horizontal trace, that naturally corresponds to \mathcal{C} -diagrams on the cylinder.

Example

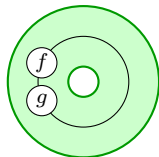
Recall that $\text{Aff}(\mathit{Braid})$ is the category of **annular braids**.

However, the strings in $\text{Tr}_h(\mathit{Braid})$ can only wrap in one direction around the cylinder, and one also has closed components wrapping around the cylinder.

Vertical trace

The **vertical trace** of a \mathbb{k} -linear category \mathcal{C} is the \mathbb{k} -module

$$\mathrm{Tr}_v(\mathcal{C}) := \left(\bigoplus_{X \in \mathrm{Ob}(\mathcal{C})} \mathrm{End}_{\mathcal{C}}(X) \right) / \mathrm{Span}_{\mathbb{k}}\{f \circ g - g \circ f\},$$



where f and g range over $f: X \rightarrow Y$, $g: Y \rightarrow X$ in \mathcal{C} .

Toroidal diagrams

If \mathcal{C} is strict pivotal, $\mathrm{Tr}_v(\mathrm{Aff}(\mathcal{C}))$ corresponds to \mathcal{C} -diagrams on the **torus**.

If $\mathcal{C} \rightarrow \mathcal{M}$ is a monoidal functor, then the $\mathrm{Tr}_v(\mathrm{Aff}(\mathcal{C}))$ **acts** on $\mathrm{Tr}_v(\mathcal{M})$.
(Nesting of annulus inside torus.)

If \mathcal{C} is right or left rigid, then $\mathrm{Tr}_v(\mathcal{C}) \cong Z(\mathrm{Aff}(\mathcal{C})) := \mathrm{End}_{\mathrm{Aff}(\mathcal{C})}(\mathbf{1})$.

Final example: HOMFLYPT skein category

The category \mathcal{FOT} of **framed oriented tangles** over the disc is isomorphic to the strict monoidal category generated by objects \uparrow, \downarrow , and morphisms

$$\nearrow, \nwarrow, \searrow, \swarrow, \downarrow, \uparrow, \downarrow, \uparrow, \cup, \cap, \cup, \cap,$$

subject to the relations

$$\begin{aligned} \cup &= | = \cap, & \cap &= \cap', & \cup &= \cup', \\ \downarrow &= \downarrow, & \downarrow &= || = \downarrow, & \downarrow &= \downarrow, \end{aligned}$$

for all orientations of the strands.

Final example: HOMFLYPT skein category

Work over $\mathbb{k} = \mathbb{Z}[z, z^{-1}, t, t^{-1}]$. The **HOMFLYPT skein category** $\mathcal{OS}(z, t)$ over the disc is the category obtained from $\mathcal{FOT}_{\mathbb{k}}$ by imposing the relations

$$\begin{array}{c} \nearrow \nearrow \\ \nwarrow \nwarrow \end{array} - \begin{array}{c} \nearrow \nwarrow \\ \nwarrow \nearrow \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array} \quad (\text{Conway skein relation})$$

and

$$\begin{array}{c} \uparrow \\ \circlearrowleft \end{array} = t \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \circlearrowright \end{array} = \frac{t - t^{-1}}{z} \mathbf{1}_{\mathbb{1}}.$$

Given an oriented link diagram L , we have

$$t^{-\text{writhe}(L)} L = H(L) \frac{t - t^{-1}}{z} \mathbf{1}_{\mathbb{1}},$$

where

- $\text{writhe}(L)$ is the **writhe number** ($\#$ pos crossings $-$ $\#$ neg crossings),
- $H(L)$ is the **HOMFLYPT** polynomial of L .

Final example: HOMFLYPT skein category

The **affinization** $\text{Aff}(\mathcal{OS}(z, t))$ is isomorphic to the **HOMFLPYT skein category** over the annulus.

Dot presentation matches the **affine oriented skein category** of Brundan.

Affinization $\text{Aff}(\mathcal{OS}(z, t))$ is obtained from $\mathcal{OS}(z, t)$ by adjoining morphisms

$$\uparrow\circlearrowleft: \uparrow \rightarrow \uparrow, \quad \downarrow\circlearrowright: \downarrow \rightarrow \downarrow,$$

subject to the relations

$$\begin{array}{c} \nearrow \\ \circlearrowleft \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \circlearrowleft \\ \nearrow \end{array}, \quad \begin{array}{c} \circlearrowleft \\ \downarrow \end{array} = \begin{array}{c} \cap \\ \downarrow\circlearrowright \end{array}, \quad \begin{array}{c} \circlearrowright \\ \downarrow \end{array} = \begin{array}{c} \cup \\ \downarrow\circlearrowleft \end{array}, \quad \uparrow\circlearrowleft \text{ is invertible.}$$

The **vertical trace** of the affinization

$$\text{Tr}_v(\text{Aff}(\mathcal{OS}(z, t)))$$

is the **HOMFLYPT skein algebra of the torus**. It acts on $\text{Tr}_v(\mathcal{OS}(z, t))$, which is the **skein of the solid torus**. (See work of Morton–Samuelson.)