

Quantum Heisenberg categorification

$$\begin{array}{c} \nearrow \\ \searrow \\ \swarrow \\ \nwarrow \end{array} - \begin{array}{c} \nwarrow \\ \swarrow \\ \nearrow \\ \searrow \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array}$$
$$\begin{array}{c} \nearrow \\ \searrow \\ \swarrow \\ \nwarrow \end{array} \text{ (red circle on bottom-right strand)} = \begin{array}{c} \nwarrow \\ \swarrow \\ \nearrow \\ \searrow \end{array} \text{ (red circle on top-right strand)}$$

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Slides available online: alistairsavage.ca/talks

Preprint: [1812.04779](https://arxiv.org/abs/1812.04779) (with J. Brundan and B. Webster)

Outline

Goal:

- 1 Define a family of quantum Heisenberg categories categorifying the Heisenberg algebra
- 2 Study categorical actions and applications in representation theory

Overview:

- 1 Strict monoidal categories and string diagrams
- 2 Quantum Heisenberg category
- 3 Categorical actions
- 4 Future directions

Strict monoidal categories

A **strict monoidal category** is a category \mathcal{C} equipped with

- a bifunctor (the **tensor product**) $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and
- a **unit object** $\mathbb{1}$,

such that, for objects A, B, C and morphisms f, g, h ,

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$,
- $(f \otimes g) \otimes h = f \otimes (g \otimes h)$,
- $1_{\mathbb{1}} \otimes f = f = f \otimes 1_{\mathbb{1}}$.

Remark: Non-strict monoidal categories

In a (not necessarily strict) **monoidal category**, the equalities above are replaced by isomorphism, and we impose some **coherence conditions**.

Every monoidal category is monoidally equivalent to a strict one.

\mathbb{k} -linear monoidal categories

Fix a commutative ground ring \mathbb{k} .

A **strict \mathbb{k} -linear monoidal category** is a strict monoidal category such that

- each morphism space is a \mathbb{k} -module,
- composition of morphisms is \mathbb{k} -bilinear,
- tensor product of morphisms is \mathbb{k} -bilinear.

The interchange law

The axioms of a strict monoidal category imply the **interchange law**: For $A_1 \xrightarrow{f} A_2$ and $B_1 \xrightarrow{g} B_2$, the following diagram commutes:

$$\begin{array}{ccc} A_1 \otimes B_1 & \xrightarrow{1 \otimes g} & A_1 \otimes B_2 \\ f \otimes 1 \downarrow & \searrow f \otimes g & \downarrow f \otimes 1 \\ A_2 \otimes B_1 & \xrightarrow{1 \otimes g} & A_2 \otimes B_2 \end{array}$$

Strict monoidal categories

Example (Monoids)

A (strict) monoidal category with one object is simply a commutative monoid. More precisely, the endomorphisms of $\mathbb{1}$ form a commutative monoid.

Conversely, every commutative monoid gives rise to a one-object monoidal category.

Example (Associative algebras)

A (strict) \mathbb{k} -linear monoidal category with one object is simply a commutative associative unital \mathbb{k} -algebra.

Categorification via split Grothendieck group

Suppose \mathcal{C} is an additive category (i.e. have \oplus).

$\text{Iso}_{\mathbb{Z}}(\mathcal{C}) =$ free abelian group generated by isom. classes of objects in \mathcal{C} .

The **split Grothendieck group** of \mathcal{C} is

$$K_0(\mathcal{C}) = \text{Iso}_{\mathbb{Z}}(\mathcal{C}) / \langle [X \oplus Y] = [X] + [Y] \mid X, Y \in \mathcal{C} \rangle.$$

If \mathcal{C} is **monoidal**, then $K_0(\mathcal{C})$ is a ring:

$$[X] \cdot [Y] = [X \otimes Y].$$

Categorification

For our purposes, to **categorify** a ring R is to find an additive monoidal category \mathcal{C} such that

$$K_0(\mathcal{C}) \cong R \quad \text{as rings.}$$

The Heisenberg algebra

Fix a **central charge** $k \in \mathbb{Z}$.

Definition

The **rank one Heisenberg algebra** has generators p^+ and p^- and relation

$$p^+p^- = p^-p^+ + k.$$

This is called the **canonical commutation relation** in physics.

Definition

The **infinite rank Heisenberg algebra** Heis_k has generators p_n^\pm , $n \geq 1$, and relations

$$p_m^+p_n^+ = p_n^+p_m^+, \quad p_m^-p_n^- = p_n^-p_m^-, \quad p_m^+p_n^- = p_n^-p_m^+ + \delta_{n,m}kn.$$

Plays a fundamental role in QFT and the theory of affine Lie algebras.

Goal: Categorify the infinite rank Heisenberg algebra.

String diagrams

Let's draw pictures! Fix a strict monoidal category \mathcal{C} .

We will denote a morphism $f: A \rightarrow B$ by:



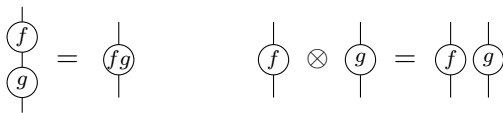
The **identity map** $1_A: A \rightarrow A$ is a string with no label:



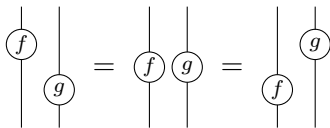
We sometimes omit the object labels when they are clear or unimportant.

String diagrams

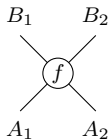
Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:



The **interchange law** then becomes:



A morphism $f: A_1 \otimes A_2 \rightarrow B_1 \otimes B_2$ can be depicted:



Presentations of strict monoidal categories

One can give **presentations** of some strict \mathbb{k} -linear monoidal categories, just as for monoids, groups, algebras, etc.

Objects: If the objects are generated by some collection $A_i, i \in I$, then we have all possible tensor products of these objects:

$$\mathbb{1}, \quad A_i, \quad A_i \otimes A_j \otimes A_k \otimes A_\ell, \quad \text{etc.}$$

Morphisms: If the morphisms are generated by some collection $f_j, j \in J$, then we have all possible compositions and tensor products of these morphisms (whenever these make sense):

$$1_{A_i}, \quad f_j \otimes (f_i f_k) \otimes (f_\ell), \quad \text{etc.}$$

We then often impose some **relations** on these morphism spaces.

String diagrams: We can build complex diagrams out of our simple generating diagrams.

Monoidally generated symmetric groups

Define a strict \mathbb{k} -linear monoidal category \mathcal{S} with one generating object \uparrow and denote

$$1_{\uparrow} = \uparrow$$

We have one generating morphism

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow.$$

We impose the relations:

$$\begin{array}{c} \nearrow \\ \cup \\ \searrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \times \\ \searrow \end{array}.$$

It is straightforward to verify that

$$\text{End}_{\mathcal{S}}(\uparrow^{\otimes n}) = \mathbb{k}S_n$$

is the group algebra of the **symmetric group** on n letters.

Monoidally generated symmetric groups

This monoidal presentation of $\mathbb{k}S_n$ is very efficient! We only needed

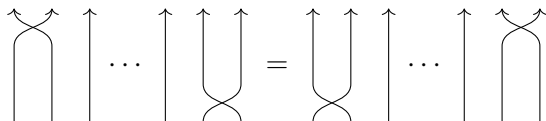
- one generating morphism, and
- two relations,

to get **all** the symmetric groups.

Note that the “distant braid relation”

$$s_i s_j = s_j s_i, \quad |i - j| > 1$$

for simple transpositions follows for free from the interchange law:



Monoidally generated braid groups

Consider the strict \mathbb{k} -linear monoidal category \mathcal{Braid} generated by

- one object \uparrow , and
- two morphisms

$$\begin{array}{c} \nearrow \\ \searrow \end{array}, \begin{array}{c} \nearrow \\ \nearrow \\ \searrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow.$$

subject to the relations

$$\begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \searrow \\ \nearrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \searrow \\ \searrow \\ \nearrow \\ \nearrow \\ \searrow \end{array}.$$

Then

$$\text{End}_{\mathcal{Braid}}(\uparrow^{\otimes n})$$

is the group algebra of the **braid group** of type A_{n-1} .

Again, this presentation is very efficient.

Monoidally generated Hecke algebras

Fix a parameter $z \in \mathbb{k}^\times$.

Let $\mathcal{H}(z)$ be the strict \mathbb{k} -linear monoidal category obtained from *Braid* by imposing one additional relation:

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} - \begin{array}{c} \nearrow \\ \times \\ \swarrow \end{array} = z \begin{array}{c} \uparrow \\ | \\ \uparrow \end{array} \quad (\text{skein relation})$$

Then

$$\text{End}_{\mathcal{H}(z)}(\uparrow^{\otimes n})$$

is the **Iwahori–Hecke algebra** of type A_{n-1} .

Monoidally generated affine Hecke algebras

Let $\mathcal{AH}(z)$ be the strict \mathbb{k} -linear monoidal category obtained from $\mathcal{H}(z)$ by adding one more **invertible** generator:

$$\uparrow \circlearrowleft : \uparrow \rightarrow \uparrow$$

and one additional relation:

$$\begin{array}{c} \nearrow \\ \searrow \\ \circlearrowleft \end{array} = \begin{array}{c} \circlearrowleft \\ \nearrow \\ \searrow \end{array} .$$

Then

$$\text{End}_{\mathcal{AH}(z)}(\uparrow^{\otimes n})$$

is the **affine Hecke algebra** of type A_{n-1} .

Algebraic versus diagrammatic

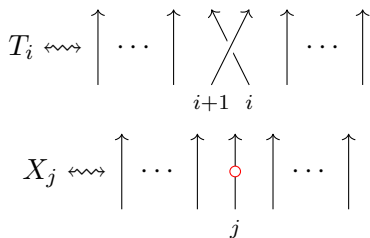
Let's match the diagrammatic description with the usual algebraic definition of the affine Hecke algebra.

Algebraic generators:

- $T_i, 1 \leq i \leq n - 1$
- $X_j^{\pm 1}, 1 \leq j \leq n$

Conventions:

- Number strands right to left
- Read diagrams bottom to top



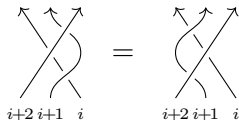
Algebraic versus diagrammatic

Relations:

① $T_i T_j = T_j T_i$ when $|i - j| > 1$

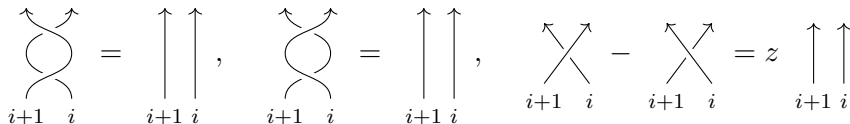
Follows from interchange relation

② $T_i T_{i+1} T_i = T_{i+1} T_i T_{i+1}$



③ $T_i^2 = z T_i + 1$

Equivalent to: T_i is invertible and $T_i - T_i^{-1} = z$

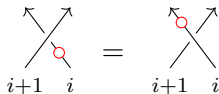


Algebraic versus diagrammatic

4 $T_i X_j = X_j T_i$ when $j \neq i, i + 1$

Follows from the interchange relation

5 $T_i X_i = X_{i+1} T_i^{-1}$



Duals

Suppose a strict monoidal category \mathcal{C} has two objects \uparrow and \downarrow , with

$$1_{\uparrow} = \uparrow \quad , \quad 1_{\downarrow} = \downarrow .$$

A morphism $\mathbb{1} \rightarrow \downarrow \otimes \uparrow$ would have string diagram

$$\begin{array}{c} \swarrow \quad \searrow \\ \downarrow \\ \vdots \end{array} \quad , \quad \text{where} \quad \begin{array}{c} \vdots \\ \vdots \end{array} = 1_{\mathbb{1}} .$$

We typically omit the dotted line and draw:

$$\begin{array}{c} \swarrow \quad \searrow \\ \cup \end{array} : \mathbb{1} \rightarrow \downarrow \otimes \uparrow .$$

Similarly, we can have

$$\begin{array}{c} \swarrow \quad \searrow \\ \cap \end{array} : \uparrow \otimes \downarrow \rightarrow \mathbb{1} .$$

Duals

We say that \downarrow is **right dual** to \uparrow (and \uparrow is **left dual** to \downarrow) if there exist morphisms

$$\cup : \mathbb{1} \rightarrow \downarrow \otimes \uparrow \quad \text{and} \quad \cap : \uparrow \otimes \downarrow \rightarrow \mathbb{1}.$$

such that

$$\downarrow \cup = \downarrow \quad \text{and} \quad \cap \uparrow = \uparrow.$$

(Analogous to the unit-counit formulation of adjunction of functors.)

The quantum Heisenberg category

Recall the **canonical commutation relation** of the Heisenberg algebra:

$$p^- p^+ = p^+ p^- - k.$$

(Assume $k < 0$ for simplicity.)

We want to **categorify** this relation. So we want objects \uparrow and \downarrow such that

$$\downarrow \otimes \uparrow \cong \uparrow \otimes \downarrow \oplus \mathbb{1}^{\oplus(-k)}. \quad (\text{Canonical commutation isom})$$

Then, in the Grothendieck ring, we have

$$[\downarrow][\uparrow] = [\uparrow][\downarrow] - k.$$

This is the canonical commutation relation with

$$[\uparrow] \longleftrightarrow p^+ \quad \text{and} \quad [\downarrow] \longleftrightarrow p^-.$$

The quantum Heisenberg category

To obtain the **quantum Heisenberg category** $\mathcal{H}eis_k(z, t)$ from $\mathcal{AH}(z)$ we perform two steps:

- 1 We adjoin a right dual \downarrow to \uparrow . Precisely, we add a generating object \downarrow and additional generating morphisms

$$\begin{array}{c} \curvearrowright \\ \uparrow \end{array} : \mathbb{1} \rightarrow \downarrow \otimes \uparrow \quad \text{and} \quad \begin{array}{c} \curvearrowleft \\ \downarrow \end{array} : \uparrow \otimes \downarrow \rightarrow \mathbb{1}$$

such that

$$\begin{array}{c} \downarrow \\ \downarrow \end{array} = \downarrow \quad \text{and} \quad \begin{array}{c} \uparrow \\ \uparrow \end{array} = \uparrow.$$

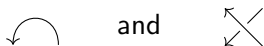
- 2 We add more generating morphisms and relations ensuring that the resulting monoidal category is pivotal and that

$$\downarrow \otimes \uparrow \cong \uparrow \otimes \downarrow \oplus \mathbb{1}^{\oplus(-k)}. \quad (\text{Canonical commutation isom})$$

The quantum Heisenberg category

There are three equivalent ways to do this. For simplicity, suppose $k = -1$.

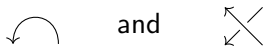
First approach: Add generating morphisms



and relations

$$\left[\begin{array}{c} \text{crossing} \\ \text{cup} \\ -tz \end{array} \right] = \left[\text{cup} \text{ crossing} \text{ cup} \right]^{-1} \quad \text{and} \quad \text{circle with dot} = tz^{-1}1_{\mathbb{1}}.$$

Second approach: Add generating morphisms

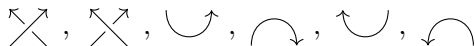


and relations

$$\left[\begin{array}{c} \text{crossing} \\ \text{cup} \\ t^{-1}z \end{array} \right] = \left[\text{cup} \text{ crossing} \text{ cup} \right]^{-1} \quad \text{and} \quad \text{circle with dot} = -t^{-1}z^{-1}1_{\mathbb{1}}.$$

The quantum Heiseberg category

Third approach: Generating morphisms



subject to the relations

$$\begin{array}{c} \text{cup} \\ \text{cap} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cap} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cross} \\ \text{cross} \end{array} = \begin{array}{c} \text{cross} \\ \text{cross} \end{array},$$

$$\begin{array}{c} \text{cross} \\ \text{cross} \end{array} - \begin{array}{c} \text{cross} \\ \text{cross} \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \downarrow, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \uparrow,$$

$$\begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \begin{array}{c} \uparrow \\ \downarrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array} + tz \begin{array}{c} \text{cup} \\ \text{cup} \end{array}, \quad \text{cap} = 0, \quad \text{cap} = -t^{-1}z^{-1}1_{\mathbb{1}},$$

and one more relation (\dagger).

The quantum Heisenberg category

Third approach: Note that we do not need the dot generator. It can be recovered via

$$\begin{array}{c} \uparrow \\ | \\ \circ \end{array} = t \begin{array}{c} \uparrow \\ | \\ \circlearrowleft \end{array} - t^2 \begin{array}{c} \uparrow \\ | \\ \end{array} .$$

The extra relation (\dagger) is that this dot is invertible.

Theorem (Brundan–S.–Webster)

- 1 All three approaches define isomorphic categories ($\mathcal{H}eis_k(z, t)$).
- 2 $\mathcal{H}eis_k(z, t)$ is **strictly pivotal** (i.e. we have isotopy invariance for morphisms).

Special cases

Deformed Heisenberg category ($k = -1$)

$\mathcal{H}eis_{-1}(z, t)$ is closely related to a **deformed Heisenberg category** $\mathcal{H}(q^2)$ introduced by Licata–S. (2013).

Precisely, $\mathcal{H}(q^2)$ is the monoidal subcategory of

$$\mathcal{H}eis_{-1}(z, -z^{-1}), \quad z = q - q^{-1},$$

consisting of all objects and morphisms that **do not involve negative powers of the dots**.

Affine oriented skein category ($k = 0$)

$\mathcal{H}eis_0(z, t)$ is the **affine oriented skein category**, an affinization of the HOMFLY-PT skein category.

Heisenberg categorification

We want an algebra isomorphism

$$K_0(\mathcal{H}eis_k(z, t)) \cong \text{Heis}_k.$$

Under this isomorphism, we should have

$$[\uparrow] \mapsto p_1^+, \quad [\downarrow] \mapsto p_1^-.$$

Questions/problems

- 1 $K_0(\mathcal{H}eis_k(z, t))$ doesn't make sense! $\mathcal{H}eis_k(z, t)$ is not an additive category!
- 2 What about the other generators p_n^\pm , $n > 1$?

Answers

- 1 Formally add in direct sums.
- 2 Enlarge the category by taking the **idempotent completion**.

Additive envelope

Suppose \mathcal{C} is some \mathbb{k} -linear monoidal category.

Its **additive envelope** is the category whose:

- **objects** are formal finite direct sums $\bigoplus_i X_i$ of objects X_i in \mathcal{C} ,
- **morphisms**

$$f: \bigoplus_{i=1}^n X_i \rightarrow \bigoplus_{j=1}^m Y_j$$

are $m \times n$ matrices, where the (j, i) -entry is a morphism

$$f_{i,j}: X_i \rightarrow Y_j.$$

Composition is given by matrix multiplication.

Idempotent completion

Definition: idempotent completion

The **idempotent completion** of a category \mathcal{C} is the category whose

- objects are pairs (A, e) where $A \in \text{Ob } \mathcal{C}$ and $e \in \text{Mor}_{\mathcal{C}}(A, A)$ is an idempotent ($e^2 = e$), and
- morphisms from (A, e) to (B, f) are elements of $f \text{Mor}_{\mathcal{C}}(A, B)e$.

Intuition: One thinks of passing to the idempotent completion as adding in objects such that the idempotents correspond to projections onto direct summands.

$$\begin{array}{c} \begin{array}{ccccccc} & & & e & & & \\ & & & \curvearrowright & & & \\ A \cong X \oplus Y & \longrightarrow & X^{\mathcal{C}} & \longrightarrow & X \oplus Y \cong A & & \end{array} \end{array}$$

Idempotent completion

Recall that we have a natural homomorphism

$$\text{Hecke algebra} \rightarrow \text{End}(\uparrow^{\otimes n}).$$

Result: Idempotents in the Hecke algebra yield idempotents in $\text{End}(\uparrow^{\otimes n})$.

The idempotents in the Hecke algebra are well known. These give us additional objects in the idempotent completion corresponding to the p_n^\pm .

For the experts (knowing about symmetric functions)

- To each partition λ of n , we have a **Young idempotent** e_λ .
- The object $(\uparrow^{\oplus n}, e_\lambda)$ corresponds to the **Schur function** s_λ .
- The p_n^\pm are the **power sums**.

Categorification of the Heisenberg algebra

There exists an injective algebra homomorphism

$$\text{Heis}_k \rightarrow K_0(\text{Kar}(\mathcal{H}eis_k(z, t))).$$

Conjecture: This is an isomorphism.

Difficulty: Passage to the Karoubi envelope. How do we know we've found all idempotents?

Degenerate case

- Previously, Khovanov defined a Heisenberg category which is essentially the $k = -1$, $z = 0$ version of the category defined here.
- Khovanov also conjectured the isomorphism.
- Then degenerate setting extended to general k (Mackaay-S., Brundan).
- Categorification conjecture recently proved in the degenerate setting (Brundan-Webster-S.).

Categorical actions ($k \neq 0$)

When $k \neq 0$, the category $\mathcal{H}eis_k(z, t)$ acts naturally on modules for **cyclotomic Hecke algebras** H_n^f of level $|k|$.

We have a chain of algebras

$$\mathbb{k} = H_0^f \subseteq H_1^f \subseteq H_2^f \subseteq \dots$$

If $k < 0$, then

- \uparrow acts by induction from $H_n^f\text{-mod}$ to $H_{n+1}^f\text{-mod}$,
- \downarrow acts by restriction from $H_n^f\text{-mod}$ to $H_{n-1}^f\text{-mod}$.

The morphisms (diagrams) act by certain natural transformations.

Fact that $\mathcal{H}eis_k(z, t)$ is pivotal corresponds to fact that induction and restriction are biadjoint.

In other words H_n^f is a **Frobenius extension** of H_{n-1}^f .

Categorical actions ($k = 0$)

Suppose $k = 0$ and $t = q^n$.

$\mathcal{H}eis_0(z, t)$ acts on representations of $U_q(\mathfrak{gl}_n)$:

- \uparrow tensors with natural module V ,
- \downarrow tensors with dual V^* .

This action extends the monoidal functor

HOMFLY-PT skein category \rightarrow cat of fd $U_q(\mathfrak{gl}_n)$ -modules

originally constructed by Turaev.

The center of $\mathcal{H}eis_0(z, t)$ maps surjectively to the center of $U_q(\mathfrak{gl}_n)$. So we get a diagrammatic calculus for this center.

Basis theorem

For many applications, one needs to know a **basis for morphism spaces** in $\mathcal{H}eis_k(z, t)$.

Usual approach: Use

- categorical actions described above,
- known bases for the algebras involved (H_n^f and $U_q(\mathfrak{gl}_n)$),
- asymptotic faithfulness (as $n \rightarrow \infty$).

However, this approach fails for $\mathcal{H}eis_k(z, t)$, $k \neq 0$, due to $\mathcal{H}eis_k(z, t)$ having a larger center than expected.

Solution: Define a categorical coproduct (“unfurling” technique of B. Webster) to define a large representation.

Quantum Frobenius Heisenberg category

Generally, can incorporate a Frobenius superalgebra F to get a more general **quantum Frobenius Heisenberg category**.

Strand can now carry tokens:

$$\uparrow \bullet f \quad , \quad f \in F.$$

We have additional/modified relations:

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} - \begin{array}{c} \nearrow \\ \times \\ \swarrow \end{array} = z \sum_{b \in B} b \uparrow \bullet b^\vee \quad , \quad \text{(new skein relation)}$$

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \\ \bullet f \end{array} = \begin{array}{c} \nearrow \\ \bullet f \\ \times \\ \searrow \end{array} \quad , \quad \begin{array}{c} \nearrow \\ \times \\ \searrow \\ \bullet f \end{array} = \begin{array}{c} \nearrow \\ \bullet f \\ \times \\ \swarrow \end{array} \quad ,$$

$$\begin{array}{c} \uparrow \\ \bullet f \\ \circ \end{array} = \begin{array}{c} \uparrow \\ \circ \\ \bullet f \end{array} \quad ,$$

+ inversion, etc.

Categorical actions

Categorical actions: Largely unexplored.

Case: $k = 0$

Obtain a **Frobenius deformation** of the affine oriented skein category.

Natural action is an open question for general F .

Should act on modules for some F -deformation of $U_q(\mathfrak{gl}_n)$.

Case: $k \neq 0$

Acts on cyclotomic quotients of **quantum affine wreath algebras** (generalize affine Hecke algebras, affine Sergeev algebras, etc.).

The structure theory of these algebras has been explored with D. Rosso.

Future directions

Traces

The **trace** of the deformed Heisenberg category $\mathcal{H}(q^2)$ was computed by Cautis–Lauda–Licata–Samuelson–Sussan.

It is related to the **elliptic Hall algebra**.

One should be able to extend this description to the larger quantum Heisenberg category $\mathcal{Heis}_k(z, t)$.

Connections to Kac–Moody 2-categories (with Brundan & Webster)

Given certain categorical Heisenberg actions, one can define a categorical Kac–Moody action.

Conversely, given certain categorical Kac–Moody actions, one can define a categorical Heisenberg action.

This extends work with Queffelec and Yacobi, which considered the level one case.