

# Quantum Heisenberg categorification

$$\begin{array}{c} \nearrow \\ \nwarrow \\ \swarrow \\ \searrow \end{array} - \begin{array}{c} \nwarrow \\ \nearrow \\ \swarrow \\ \searrow \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array} \qquad \begin{array}{c} \nwarrow \\ \nearrow \\ \swarrow \\ \searrow \end{array} \text{ (with red dot on } \swarrow \text{)} = \begin{array}{c} \nwarrow \\ \nearrow \\ \swarrow \\ \searrow \end{array} \text{ (with red dot on } \nearrow \text{)}$$

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Slides available online: [alistairsavage.ca/talks](http://alistairsavage.ca/talks)

Preprint: [1812.04779](https://arxiv.org/abs/1812.04779) (with J. Brundan and B. Webster)

# Outline

## Goal:

- 1 Define a family of quantum Heisenberg categories categorifying the Heisenberg algebra
- 2 Study categorical actions and applications in representation theory

## Overview:

- 1 Strict monoidal categories and string diagrams
- 2 Quantum Heisenberg category
- 3 Categorical actions
- 4 Future directions

# Strict monoidal categories

A **strict monoidal category** is a category  $\mathcal{C}$  equipped with

- a bifunctor (the **tensor product**)  $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ , and
- a **unit object**  $\mathbb{1}$ ,

such that

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$  for all objects  $A, B, C$ ,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$  for all objects  $A$ .

## Remark: Non-strict monoidal categories

In a (not necessarily strict) **monoidal category**, the equalities above are replaced by isomorphism, and we impose some **coherence conditions**.

Every monoidal category is monoidally equivalent to a strict one.

## $\mathbb{k}$ -linear monoidal categories

Fix a commutative ground ring  $\mathbb{k}$ .

A **strict  $\mathbb{k}$ -linear monoidal category** is a strict monoidal category such that

- each morphism space is a  $\mathbb{k}$ -module,
- composition of morphisms is  $\mathbb{k}$ -bilinear,
- tensor product of morphisms is  $\mathbb{k}$ -bilinear.

### The interchange law

The axioms of a strict monoidal category imply the **interchange law**: For  $A_1 \xrightarrow{f} A_2$  and  $B_1 \xrightarrow{g} B_2$ , the following diagram commutes:

$$\begin{array}{ccc} A_1 \otimes B_1 & \xrightarrow{1 \otimes g} & A_1 \otimes B_2 \\ f \otimes 1 \downarrow & \searrow f \otimes g & \downarrow f \otimes 1 \\ A_2 \otimes B_1 & \xrightarrow{1 \otimes g} & A_2 \otimes B_2 \end{array}$$

# Strict monoidal categories

## Example (Monoids)

A (strict) monoidal category with one object is simply a commutative monoid. More precisely, the endomorphisms of  $\mathbb{1}$  form a commutative monoid.

Conversely, every commutative monoid gives rise to a one-object monoidal category.

## Example (Associative algebras)

A (strict)  $\mathbb{k}$ -linear monoidal category is simply a commutative associative unital  $\mathbb{k}$ -algebra.

## Categorification via split Grothendieck group

Suppose  $\mathcal{C}$  is an additive category (i.e. have  $\oplus$ ).

$\text{Iso}_{\mathbb{Z}}(\mathcal{C}) =$  free abelian group generated by isom. classes of objects in  $\mathcal{C}$ .

The **split Grothendieck group** of  $\mathcal{C}$  is

$$K_0(\mathcal{C}) = \text{Iso}_{\mathbb{Z}}(\mathcal{C}) / \langle [X \oplus Y] = [X] + [Y] \mid X, Y \in \mathcal{C} \rangle.$$

If  $\mathcal{C}$  is **monoidal**, then  $K_0(\mathcal{C})$  is a ring:

$$[X] \cdot [Y] = [X \otimes Y].$$

### Categorification

For our purposes, to **categorify** a ring  $R$  is to find an additive monoidal category  $\mathcal{C}$  such that

$$K_0(\mathcal{C}) \cong R \quad \text{as rings.}$$

# The Heisenberg algebra

Fix a **central charge**  $k \in \mathbb{Z}$ .

## Definition

The **rank one Heisenberg algebra** has generators  $p^+$  and  $p^-$  and relation

$$p^+p^- = p^-p^+ + k.$$

This is called the **canonical commutation relation** in physics.

## Definition

The **infinite rank Heisenberg algebra**  $\text{Heis}_k$  has generators  $p_n^\pm$ ,  $n \geq 1$ , and relations

$$p_m^+p_n^+ = p_n^+p_m^+, \quad p_m^-p_n^- = p_n^-p_m^-, \quad p_m^+p_n^- = p_n^-p_m^+ + \delta_{n,m}kn.$$

Plays a fundamental role in QFT and the theory of affine Lie algebras.

**Goal:** Categorify the infinite rank Heisenberg algebra.

## String diagrams

Let's draw pictures! Fix a strict monoidal category  $\mathcal{C}$ .

We will denote a morphism  $f: A \rightarrow B$  by:



The **identity map**  $1_A: A \rightarrow A$  is a string with no label:

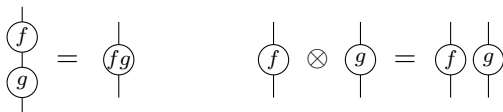


We sometimes omit the object labels when they are clear or unimportant.

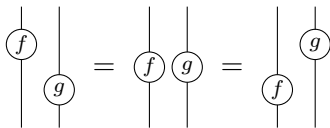


# String diagrams

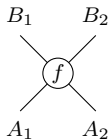
Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:



The **interchange law** then becomes:



A morphism  $f: A_1 \otimes A_2 \rightarrow B_1 \otimes B_2$  can be depicted:



# Presentations of strict monoidal categories

One can give **presentations** of some strict  $\mathbb{k}$ -linear monoidal categories, just as for monoids, groups, algebras, etc.

**Objects:** If the objects are generated by some collection  $A_i, i \in I$ , then we have all possible tensor products of these objects:

$$\mathbb{1}, \quad A_i, \quad A_i \otimes A_j \otimes A_k \otimes A_\ell, \quad \text{etc.}$$

**Morphisms:** If the morphisms are generated by some collection  $f_j, j \in J$ , then we have all possible compositions and tensor products of these morphisms (whenever these make sense):

$$1_{A_i}, \quad f_j \otimes (f_i f_k) \otimes (f_\ell), \quad \text{etc.}$$

We then often impose some **relations** on these morphism spaces.

**String diagrams:** We can build complex diagrams out of our simple generating diagrams.

# Monoidally generated symmetric groups

Define a strict  $\mathbb{k}$ -linear monoidal category  $\mathcal{S}$  with one generating object  $\uparrow$  and denote

$$1_{\uparrow} = \uparrow$$

We have one generating morphism

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow .$$

We impose the relations:

$$\begin{array}{c} \nearrow \\ \cup \\ \searrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} .$$

It is straightforward to verify that

$$\text{End}_{\mathcal{S}}(\uparrow^{\otimes n}) = \mathbb{k}S_n$$

is the group algebra of the **symmetric group** on  $n$  letters.

# Monoidally generated symmetric groups

This monoidal presentation of  $\mathbb{k}S_n$  is very efficient! We only needed

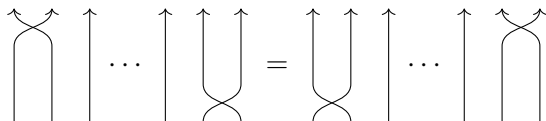
- one generating morphism, and
- two relations,

to get **all** the symmetric groups.

Note that the “distant braid relation”

$$s_i s_j = s_j s_i, \quad |i - j| > 1$$

for simple transpositions follows for free from the interchange law:



# Monoidally generated braid groups

Consider the strict  $\mathbb{k}$ -linear monoidal category  $\mathcal{Braid}$  generated by

- one object  $\uparrow$ , and
- two morphisms

$$\begin{array}{c} \nearrow \\ \searrow \end{array}, \begin{array}{c} \nearrow \\ \nearrow \\ \searrow \end{array} : \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow.$$

subject to the relations

$$\begin{array}{c} \uparrow \\ \curvearrowright \\ \uparrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \uparrow \\ \curvearrowleft \\ \uparrow \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} = \begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array}.$$

Then

$$\text{End}_{\mathcal{Braid}}(\uparrow^{\otimes n})$$

is the group algebra of the **braid group** of type  $A_{n-1}$ .

Again, this presentation is very efficient.

# Monoidally generated Hecke algebras

Fix a parameter  $z \in \mathbb{k}^\times$ .

Let  $\mathcal{H}(z)$  be the strict  $\mathbb{k}$ -linear monoidal category obtained from *Braid* by imposing one additional relation:

$$\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} - \begin{array}{c} \nearrow \\ \times \\ \nearrow \end{array} = z \begin{array}{c} \uparrow \\ | \\ \uparrow \end{array} \begin{array}{c} \uparrow \\ | \\ \uparrow \end{array} \quad (\text{skein relation})$$

Then

$$\text{End}_{\mathcal{H}(z)}(\uparrow^{\otimes n})$$

is the **Iwahori–Hecke algebra** of type  $A_{n-1}$ .

# Monoidally generated affine Hecke algebras

Let  $\mathcal{AH}(z)$  be the strict  $\mathbb{k}$ -linear monoidal category obtained from  $\mathcal{H}(z)$  by adding one more **invertible** generator:

$$\uparrow \circlearrowleft : \uparrow \rightarrow \uparrow$$

and one additional relation:

$$\begin{array}{c} \nearrow \\ \searrow \\ \circlearrowleft \end{array} = \begin{array}{c} \circlearrowleft \\ \nearrow \\ \searrow \end{array} .$$

Then

$$\text{End}_{\mathcal{AH}(z)}(\uparrow^{\otimes n})$$

is the **affine Hecke algebra** of type  $A_{n-1}$ .

# Algebraic versus diagrammatic

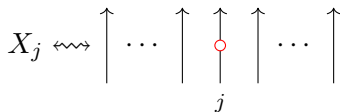
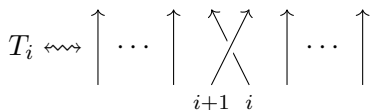
Let's match the diagrammatic description with the usual algebraic definition of the affine Hecke algebra.

Algebraic generators:

- $T_i, 1 \leq i \leq n - 1$
- $X_j^{\pm 1}, 1 \leq j \leq n$

Conventions:

- Number strands right to left
- Read diagrams bottom to top





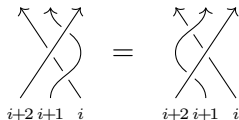
# Algebraic versus diagrammatic

Relations:

①  $T_i T_j = T_j T_i$  when  $|i - j| > 1$

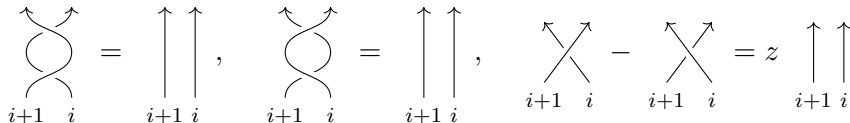
Follows from interchange relation

②  $T_i T_{i+1} T_i = T_{i+1} T_i T_{i+1}$



③  $T_i^2 = z T_i + 1$

Equivalent to:  $T_i$  is invertible and  $T_i - T_i^{-1} = z$

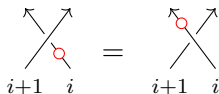


# Algebraic versus diagrammatic

④  $T_i X_j = X_j T_i$  when  $j \neq i, i + 1$

Follows from the interchange relation

⑤  $T_i X_i = X_{i+1} T_i^{-1}$



# Duals

Suppose a strict monoidal category  $\mathcal{C}$  has two objects  $\uparrow$  and  $\downarrow$ , with

$$1_{\uparrow} = \uparrow \quad , \quad 1_{\downarrow} = \downarrow .$$

A morphism  $\mathbb{1} \rightarrow \downarrow \otimes \uparrow$  would have string diagram

$$\begin{array}{c} \swarrow \quad \searrow \\ \downarrow \\ \vdots \end{array} \quad , \quad \text{where} \quad \begin{array}{c} \vdots \\ \vdots \end{array} = 1_{\mathbb{1}} .$$

We typically omit the dotted line and draw:

$$\begin{array}{c} \swarrow \quad \searrow \\ \cup \end{array} : \mathbb{1} \rightarrow \downarrow \otimes \uparrow .$$

Similarly, we can have

$$\begin{array}{c} \swarrow \quad \searrow \\ \cap \end{array} : \uparrow \otimes \downarrow \rightarrow \mathbb{1} .$$

# Duals

We say that  $\downarrow$  is **right dual** to  $\uparrow$  (and  $\uparrow$  is **left dual** to  $\downarrow$ ) if there exist morphisms

$$\cup : \mathbb{1} \rightarrow \downarrow \otimes \uparrow \quad \text{and} \quad \cap : \uparrow \otimes \downarrow \rightarrow \mathbb{1}.$$

such that

$$\downarrow \cup = \downarrow \quad \text{and} \quad \cap \uparrow = \uparrow.$$

(Analogous to the unit-counit formulation of adjunction of functors.)

# The quantum Heisenberg category

Recall the **canonical commutation relation** of the Heisenberg algebra:

$$p^- p^+ = p^+ p^- - k.$$

(Assume  $k < 0$  for simplicity.)

We want to **categorify** this relation. So we want objects  $\uparrow$  and  $\downarrow$  such that

$$\downarrow \otimes \uparrow \cong \uparrow \otimes \downarrow \oplus \mathbb{1}^{\oplus(-k)}. \quad (\text{Canonical commutation isom})$$

Then, in the Grothendieck ring, we have

$$[\downarrow][\uparrow] = [\uparrow][\downarrow] - k.$$

This is the canonical commutation relation with

$$[\uparrow] \longleftrightarrow p^+ \quad \text{and} \quad [\downarrow] \longleftrightarrow p^-.$$

# The quantum Heisenberg category

To obtain the **quantum Heisenberg category**  $\mathcal{H}eis_k(z, t)$  from  $\mathcal{AH}(z)$  we perform two steps:

- 1 We adjoin a right dual  $\downarrow$  to  $\uparrow$ . Precisely, we add a generating object  $\downarrow$  and additional generating morphisms

$$\cup : \mathbb{1} \rightarrow \downarrow \otimes \uparrow \quad \text{and} \quad \cap : \uparrow \otimes \downarrow \rightarrow \mathbb{1}$$

such that

$$\downarrow \cup \uparrow = \downarrow \quad \text{and} \quad \uparrow \cap \downarrow = \uparrow.$$

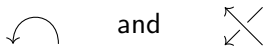
- 2 We add more generating morphisms and relations ensuring that the resulting monoidal category is pivotal and that

$$\downarrow \otimes \uparrow \cong \uparrow \otimes \downarrow \oplus \mathbb{1}^{\oplus(-k)}. \quad (\text{Canonical commutation isom})$$

# The quantum Heisenberg category

There are three equivalent ways to do this. For simplicity, suppose  $k = -1$ .

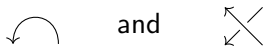
**First approach:** Add generating morphisms



and relations

$$\left[ \begin{array}{c} \text{crossing} \\ \text{cup} \\ -tz \end{array} \right] = \left[ \text{cup} \text{ crossing} \text{ cup} \right]^{-1} \quad \text{and} \quad \text{circle with dot} = tz^{-1}1_{\mathbb{1}}.$$

**Second approach:** Add generating morphisms

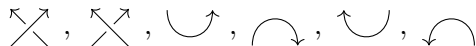


and relations

$$\left[ \begin{array}{c} \text{crossing} \\ \text{cup} \\ t^{-1}z \end{array} \right] = \left[ \text{cup} \text{ crossing} \text{ cup} \right]^{-1} \quad \text{and} \quad \text{circle with dot} = -t^{-1}z^{-1}1_{\mathbb{1}}.$$

# The quantum Heiseberg category

Third approach: Generating morphisms



subject to the relations

$$\begin{array}{c} \text{cup} \\ \text{cap} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cap} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cross} \\ \text{cross} \end{array} = \begin{array}{c} \text{cross} \\ \text{cross} \end{array},$$

$$\begin{array}{c} \text{cross} \\ \text{cross} \end{array} - \begin{array}{c} \text{cross} \\ \text{cross} \end{array} = z \begin{array}{c} \uparrow \\ \uparrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \downarrow, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \uparrow,$$

$$\begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \begin{array}{c} \uparrow \\ \downarrow \end{array}, \quad \begin{array}{c} \text{cup} \\ \text{cup} \end{array} = \begin{array}{c} \uparrow \\ \uparrow \end{array} + tz \begin{array}{c} \text{cup} \\ \text{cup} \end{array}, \quad \text{cap} = 0, \quad \text{cap} = -t^{-1}z^{-1}1_{\mathbb{1}},$$

and one more relation ( $\dagger$ ).



# The quantum Heisenberg category

**Third approach:** Note that we do not need the dot generator. It can be recovered via

$$\begin{array}{c} \uparrow \\ | \\ \circ \end{array} = t \begin{array}{c} \uparrow \\ | \\ \circlearrowleft \end{array} - t^2 \begin{array}{c} \uparrow \\ | \end{array} .$$

The extra relation ( $\dagger$ ) is that this dot is invertible.

## Theorem (Brundan–S.–Webster)

- 1 All three approaches define isomorphic categories ( $\mathcal{H}eis_k(z, t)$ ).
- 2  $\mathcal{H}eis_k(z, t)$  is **strictly pivotal** (i.e. we have isotopy invariance for morphisms).

## Special cases

### Deformed Heisenberg category ( $k = -1$ )

$\mathcal{H}eis_{-1}(z, t)$  is closely related to a **deformed Heisenberg category**  $\mathcal{H}(q^2)$  introduced by Licata–S. (2013).

Precisely,  $\mathcal{H}(q^2)$  is the monoidal subcategory of

$$\mathcal{H}eis_{-1}(z, -z^{-1}), \quad z = q - q^{-1},$$

consisting of all objects and morphisms that **do not involve negative powers of the dots**.

### Affine oriented skein category ( $k = 0$ )

$\mathcal{H}eis_0(z, t)$  is the **affine oriented skein category**, an affinization of the HOMFLY-PT skein category.

# Heisenberg categorification

We want an algebra isomorphism

$$K_0(\mathcal{H}eis_k(z, t)) \cong \text{Heis}_k.$$

Under this isomorphism, we should have

$$[\uparrow] \mapsto p_1^+, \quad [\downarrow] \mapsto p_1^-.$$

## Questions/problems

- 1  $K_0(\mathcal{H}eis_k(z, t))$  doesn't make sense!  $\mathcal{H}eis_k(z, t)$  is not an additive category!
- 2 What about the other generators  $p_n^\pm$ ,  $n > 1$ ?

## Answers

- 1 Formally add in direct sums.
- 2 Enlarge the category by taking the **idempotent completion**.

# Additive envelope

Suppose  $\mathcal{C}$  is some  $\mathbb{k}$ -linear monoidal category.

Its **additive envelope** is the category whose:

- **objects** are formal finite direct sums  $\bigoplus_i X_i$  of objects  $X_i$  in  $\mathcal{C}$ ,
- **morphisms**

$$f: \bigoplus_{i=1}^n X_i \rightarrow \bigoplus_{j=1}^m Y_j$$

are  $m \times n$  matrices, where the  $(j, i)$ -entry is a morphism

$$f_{i,j}: X_i \rightarrow Y_j.$$

Composition is given by matrix multiplication.

# Idempotent completion

## Definition: idempotent completion

The **idempotent completion** of a category  $\mathcal{C}$  is the category whose

- objects are pairs  $(A, e)$  where  $A \in \text{Ob } \mathcal{C}$  and  $e \in \text{Mor}_{\mathcal{C}}(A, A)$  is an idempotent ( $e^2 = e$ ), and
- morphisms from  $(A, e)$  to  $(B, f)$  are elements of  $f \text{Mor}_{\mathcal{C}}(A, B)e$ .

**Intuition:** One thinks of passing to the idempotent completion as adding in objects such that the idempotents correspond to projections onto direct summands.

$$\begin{array}{c} \begin{array}{ccccccc} & & & e & & & \\ & & & \curvearrowright & & & \\ A \cong X \oplus Y & \longrightarrow & X^{\mathcal{C}} & \longrightarrow & X \oplus Y & \cong & A \end{array} \end{array}$$

# Idempotent completion

Recall that we have a natural homomorphism

$$\text{Hecke algebra} \rightarrow \text{End}(\uparrow^{\otimes n}).$$

**Result:** Idempotents in the Hecke algebra yield idempotents in  $\text{End}(\uparrow^{\otimes n})$ .

The idempotents in the Hecke algebra are well known. These give us additional objects in the idempotent completion corresponding to the  $p_n^\pm$ .

For the experts (knowing about symmetric functions)

- To each partition  $\lambda$  of  $n$ , we have a **Young idempotent**  $e_\lambda$ .
- The object  $(\uparrow^{\oplus n}, e_\lambda)$  corresponds to the **Schur function**  $s_\lambda$ .
- The  $p_n^\pm$  are the **power sums**.

# Categorification of the Heisenberg algebra

There exists an injective algebra homomorphism

$$\text{Heis}_k \rightarrow K_0(\text{Kar}(\mathcal{H}eis_k(z, t))).$$

**Conjecture:** This is an isomorphism.

**Difficulty:** Passage to the Karoubi envelope. How do we know we've found all idempotents?

## Degenerate case

- Previously, Khovanov defined a Heisenberg category which is essentially the  $k = -1$ ,  $z = 0$  version of the category defined here.
- Khovanov also conjectured the isomorphism.
- Then degenerate setting extended to general  $k$  (Mackaay-S., Brundan).
- Categorification conjecture recently proved in the degenerate setting (Brundan–Webster–S.).

## Categorical actions ( $k \neq 0$ )

When  $k \neq 0$ , the category  $\mathcal{H}eis_k(z, t)$  acts naturally on modules for **cyclotomic Hecke algebras**  $H_n^f$  of level  $|k|$ .

We have a chain of algebras

$$\mathbb{k} = H_0^f \subseteq H_1^f \subseteq H_2^f \subseteq \dots$$

If  $k < 0$ , then

- $\uparrow$  acts by induction from  $H_n^f\text{-mod}$  to  $H_{n+1}^f\text{-mod}$ ,
- $\downarrow$  acts by restriction from  $H_n^f\text{-mod}$  to  $H_{n-1}^f\text{-mod}$ .

The morphisms (diagrams) act by certain natural transformations.

Fact that  $\mathcal{H}eis_k(z, t)$  is pivotal corresponds to fact that induction and restriction are biadjoint.

In other words  $H_n^f$  is a **Frobenius extension** of  $H_{n-1}^f$ .



## Categorical actions ( $k = 0$ )

Suppose  $k = 0$  and  $t = q^n$ .

$\mathcal{H}eis_0(z, t)$  acts on representations of  $U_q(\mathfrak{gl}_n)$ :

- $\uparrow$  tensors with natural module  $V$ ,
- $\downarrow$  tensors with dual  $V^*$ .

This action extends the monoidal functor

HOMFLY-PT skein category  $\rightarrow$  cat of fd  $U_q(\mathfrak{gl}_n)$ -modules

originally constructed by Turaev.

The center of  $\mathcal{H}eis_0(z, t)$  maps surjectively to the center of  $U_q(\mathfrak{gl}_n)$ . So we get a diagrammatic calculus for this center.

# Basis theorem

For many applications, one needs to know a **basis for morphism spaces** in  $\mathcal{H}eis_k(z, t)$ .

**Usual approach:** Use

- categorical actions described above,
- known bases for the algebras involved ( $H_n^f$  and  $U_q(\mathfrak{gl}_n)$ ),
- asymptotic faithfulness (as  $n \rightarrow \infty$ ).

**However**, this approach fails for  $\mathcal{H}eis_k(z, t)$ ,  $k \neq 0$ , due to  $\mathcal{H}eis_k(z, t)$  having a larger center than expected.

**Solution:** Define a categorical coproduct (“unfurling” technique of B. Webster) to define a large representation.

# Quantum Frobenius Heisenberg category

Generally, can incorporate a Frobenius superalgebra  $F$  to get a more general **quantum Frobenius Heisenberg category**.

Strand can now carry tokens:

$$\uparrow \bullet f \quad , \quad f \in F.$$

We have additional/modified relations:

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} - \begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} = z \sum_{b \in B} b \uparrow \bullet b^\vee \quad , \quad \text{(new skein relation)}$$

$$\begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} \bullet f = f \begin{array}{c} \nearrow \searrow \\ \nwarrow \nearrow \end{array} \quad , \quad \begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} \bullet f = f \begin{array}{c} \nwarrow \nearrow \\ \nearrow \searrow \end{array} \quad ,$$

$$\begin{array}{c} \uparrow \bullet f \\ \circ \end{array} = \begin{array}{c} \uparrow \circ \\ \bullet f \end{array} \quad ,$$

+ inversion, etc.

# Categorical actions

Categorical actions: Largely unexplored.

Case:  $k = 0$

Obtain a **Frobenius deformation** of the affine oriented skein category.

Natural action is an open question for general  $F$ .

Should act on modules for some  $F$ -deformation of  $U_q(\mathfrak{gl}_n)$ .

Case:  $k \neq 0$

Acts on cyclotomic quotients of **quantum affine wreath algebras** (generalize affine Hecke algebras, affine Sergeev algebras, etc.).

The structure theory of these algebras has been explored with D. Rosso.

# Future directions

## Traces

The **trace** of the deformed Heisenberg category  $\mathcal{H}(q^2)$  was computed by Cautis–Lauda–Licata–Samuelson–Sussan.

It is related to the **elliptic Hall algebra**.

One should be able to extend this description to the larger quantum Heisenberg category  $\mathcal{H}eis_k(z, t)$ .

## Connections to Kac–Moody 2-categories (with Brundan & Webster)

Given certain categorical Heisenberg actions, one can define a categorical Kac–Moody action.

Conversely, given certain categorical Kac–Moody actions, one can define a categorical Heisenberg action.

This extends work with Queffelec and Yacobi, which considered the level one case.