

Adjunction

uOttawa Math Club

$$\text{↯} = \uparrow \quad \text{↵} = \downarrow$$

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Dual spaces

Suppose V and W are finite-dimensional complex vector spaces.

The **dual** of V is the space of linear functionals on V :

$$V^* := \{f : V \rightarrow \mathbb{C} : f \text{ is linear}\}.$$

It is itself a vector space under

- **pointwise addition:**

$$(f + g)(v) := f(v) + g(v), \quad f, g \in V^*, \quad v \in V,$$

- **scalar multiplication:**

$$(\alpha f)(v) := \alpha f(v), \quad \alpha \in \mathbb{C}, \quad f \in V^*, \quad v \in V.$$

Transposes

Suppose we have a linear map

$$A: V \rightarrow W.$$

Then, given $f \in W^*$, we can produce an element of V^* :

$$\begin{array}{ccc} V & \xrightarrow{A} & W \\ & \searrow f \circ A & \downarrow f \\ & & \mathbb{C} \end{array}$$

The **transpose** of $A: V \rightarrow W$ is the linear map

$$\begin{aligned} A^T: W^* &\rightarrow V^*, \\ A^T f &:= f \circ A, \quad f \in W^*. \end{aligned}$$

Inner product spaces

An **inner product** on V is a map

$$\langle \cdot, \cdot \rangle_V: V \times V \rightarrow \mathbb{C},$$

satisfying the following axioms for all $u, v, w \in V$ and all $\alpha \in \mathbb{C}$:

- **conjugate symmetry:**

$$\langle u, v \rangle_V = \overline{\langle v, u \rangle_V}$$

- **linearity in the first argument:**

$$\begin{aligned}\langle \alpha v, w \rangle_V &= \alpha \langle v, w \rangle_V \\ \langle u + v, w \rangle_V &= \langle u, w \rangle_V + \langle v, w \rangle_V\end{aligned}$$

- **positive-definiteness:**

$$\begin{aligned}\langle v, v \rangle_V &\geq 0 \\ \langle v, v \rangle_V = 0 &\iff v = 0\end{aligned}$$

Inner product spaces

Example

If $V = \mathbb{C}^n$ (column vectors), then we have the **standard inner product**

$$\langle u, v \rangle = u^T \bar{v}.$$

Suppose V has an inner product. For each $v \in V$, we can define a linear map

$$f_v: V \rightarrow \mathbb{C}, \quad f_v(u) = \langle u, v \rangle_V, \quad u \in V.$$

So $f_v \in V^*$.

One can check that this defines an **isomorphism** of vector spaces

$$V \cong V^*, \quad v \mapsto f_v.$$

Summary: Inner products allow us to identify a f.d. vector space with its dual.

Adjoints

Suppose V and W both have inner products, and that we have a linear map

$$A: V \rightarrow W.$$

Then we have the following:

$$\begin{array}{ccc} W^* & \xrightarrow{A^T} & V^* \\ \cong \downarrow & & \downarrow \cong \\ W & \xrightarrow{A^*} & V \end{array}$$

We can complete the square to a commutative diagram. The bottom map is the **adjoint** A^* of A .

The adjoint is **uniquely determined** by the fact that

$$\langle Av, w \rangle_W = \langle v, A^*w \rangle_V, \quad v \in V, w \in W.$$

Adjoint

Example

Suppose $V = \mathbb{C}^n$ and $W = \mathbb{C}^m$, with the standard inner products.

Every linear map $A: V \rightarrow W$ corresponds to a matrix.

The defining property of the adjoint is that

$$(Av)^T \bar{w} = \langle Av, w \rangle = \langle v, A^* w \rangle = v^T \overline{A^* w}, \quad v \in V, w \in W.$$

So, for all $v \in V$ and $w \in W$, we have

$$v^T A^T \bar{w} = v^T \overline{A^* w} = v^T \overline{A^*} \bar{w}.$$

This implies that $A^T = \overline{A^*}$, which is equivalent to

$$A^* = \overline{A^T}.$$

So the adjoint A^* is the **conjugate transpose** of the matrix A .

Categories (Definition)

A **category** \mathcal{C} consists of

- a class of **objects** $\text{Ob } \mathcal{C}$,
- a class of **morphisms** $\text{hom}_{\mathcal{C}}(X, Y)$ for all $X, Y \in \text{Ob } \mathcal{C}$,

together with a **composition**

$$\text{hom}_{\mathcal{C}}(Y, Z) \times \text{hom}_{\mathcal{C}}(X, Y) \rightarrow \text{hom}_{\mathcal{C}}(X, Z), \quad (f, g) \mapsto f \circ g,$$

and an **identity morphism** $1_X \in \text{hom}_{\mathcal{C}}(X, X)$ for all objects $X \in \text{Ob } \mathcal{C}$.

The composition must be associative:

$$(f \circ g) \circ h = f \circ (g \circ h)$$

whenever $f \circ g$ and $g \circ h$ are defined.

The identity morphism has the property that

$$1_Y \circ f = f = f \circ 1_X \quad \text{for all } f \in \text{hom}_{\mathcal{C}}(X, Y).$$

Categories (Examples)

Example (Sets)

- **Objects:** sets
- **Morphisms:** set maps

Example (Vector spaces)

- **Objects:** vector spaces over a fixed field
- **Morphisms:** linear maps

Example (Groups)

- **Objects:** groups
- **Morphisms:** group homomorphisms

Categories (Examples)

Example (Rings)

- **Objects:** rings
- **Morphisms:** ring homomorphisms

Example (Topological spaces)

- **Objects:** topological spaces
- **Morphisms:** continuous maps

Other examples

- modules over a fixed ring
- smooth manifolds
- algebraic varieties
- ...

Functors

Suppose \mathcal{C} and \mathcal{D} are categories.

A **functor** F from \mathcal{C} to \mathcal{D} consists of

- a map $F: \text{Ob } \mathcal{C} \rightarrow \text{Ob } \mathcal{D}$,
- for all $X, Y \in \text{Ob } \mathcal{C}$, a map $F: \text{hom}_{\mathcal{C}}(X, Y) \rightarrow \text{hom}_{\mathcal{D}}(F(X), F(Y))$.

We require that the map on morphisms respects composition:

$$F(f \circ g) = F(f) \circ F(g),$$

whenever the composition of $f, g \in \text{hom } \mathcal{C}$ is defined.

We also require it to preserve identities:

$$F(1_X) = 1_{F(X)}.$$

Functors

Example (Forgetful functors)

We can define a functor

$$F: \text{category of groups} \rightarrow \text{category of sets}$$

as follows:

- for a group G , we define $F(G)$ to be the underlying set of G ,
- for a group homomorphism $f: G_1 \rightarrow G_2$, we define $F(f)$ to be the underlying set map.

So F just **forgets** the group structure.

There are many other examples of forgetful functors:

- category of rings \rightarrow category of abelian groups, $(R, +, \cdot) \mapsto (R, +)$
- category of rings \rightarrow category of sets
- category of vector spaces \rightarrow category of sets

Functors

Example (Double dual)

There is a functor from the category of complex vector spaces to itself that

- maps any vector space to its double dual (the dual of its dual space),
- maps any linear map to its double transpose.

Example (Fundamental group)

Suppose

- **Top** is the category of **pointed topological spaces** (topological spaces together with a distinguished point),
- **Group** is the category of groups.

We have a functor **Top** \rightarrow **Group** that maps a pointed topological space to its **fundamental group**.

Adjoint functors: general idea

Suppose \mathcal{C} and \mathcal{D} are categories. Let **Set** be the category of sets.

General philosophy

We will think of

$$\mathrm{hom}_{\mathcal{C}}(\cdot, \cdot): \mathcal{C} \times \mathcal{C} \rightarrow \mathbf{Set}$$

as a categorical analogue of an inner product on the category \mathcal{C} .

Goal

Define an appropriate notion of **adjoint functor**.

First guess

The adjoint of a functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is a functor $G: \mathcal{D} \rightarrow \mathcal{C}$ such that

$$\mathrm{hom}_{\mathcal{D}}(FX, Y) = \mathrm{hom}_{\mathcal{C}}(X, GY), \quad X \in \mathrm{Ob} \mathcal{C}, Y \in \mathrm{Ob} \mathcal{D}.$$

Problem 1: Lack of symmetry

For inner products, we had conjugate symmetry:

$$\langle u, v \rangle_V = \overline{\langle v, u \rangle_V}, \quad u, v \in V.$$

It follows that $(A^*)^* = A$.

There is no such symmetry (in general) in the setting of categories.

Solution: Break the symmetry!

Second guess

A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is **left adjoint** to a functor $G: \mathcal{D} \rightarrow \mathcal{C}$ (and G is **right adjoint** to F) if

$$\mathrm{hom}_{\mathcal{D}}(FX, Y) = \mathrm{hom}_{\mathcal{C}}(X, GY), \quad X \in \mathrm{Ob} \mathcal{C}, Y \in \mathrm{Ob} \mathcal{D}.$$

Problem 2: Equality is weird

Example

Define two groups as follows:

$$G = \{a, b\}, \quad a^2 = b^2 = 1, \quad ab = ba = b,$$
$$H = \{x, y\}, \quad x^2 = y^2 = 1, \quad xy = yx = y.$$

Are these the same group (i.e. are they **equal**)? Technically, no.

However, they are **isomorphic** via the map $a \mapsto x, b \mapsto y$.

In fact, both are isomorphic to the cyclic group of order 2.

Observation: In category theory, the concept of **isomorphism** is more fundamental than the notion of **equality**.

From equality to isomorphism

Our observation suggests that the condition

$$\text{hom}_{\mathcal{C}}(FX, Y) = \text{hom}_{\mathcal{D}}(X, GY), \quad X \in \text{Ob } \mathcal{C}, Y \in \text{Ob } \mathcal{D},$$

is not very natural.

Solution: Change from equality to isomorphism!

Third guess

A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is **left adjoint** to a functor $G: \mathcal{D} \rightarrow \mathcal{C}$ (and G is **right adjoint** to F) if

$$\text{hom}_{\mathcal{D}}(FX, Y) \cong \text{hom}_{\mathcal{C}}(X, GY), \quad X \in \text{Ob } \mathcal{C}, Y \in \text{Ob } \mathcal{D}.$$

(Here \cong is isomorphism of sets.)

Problem 3: Naturality

Generally, in mathematics, natural constructions should respect the structure that is present:

- 1 **vector space** homomorphisms should be linear,
- 2 **group** homomorphisms should commute with the group operation,
- 3 **ring** homomorphisms should commute with multiplication and addition.

Our current definition doesn't really fully respect all of the structure of the categories \mathcal{C} and \mathcal{D} coming from morphisms.

Suppose

$$f \in \text{hom}_{\mathcal{C}}(X_1, X_2), \quad g \in \text{hom}_{\mathcal{C}}(Y_1, Y_2).$$

Then we have a map

$$\begin{aligned} \mu_{f,g}: \text{hom}_{\mathcal{C}}(X_2, Y_1) &\rightarrow \text{hom}_{\mathcal{C}}(X_1, Y_2), \\ A &\mapsto g \circ A \circ f. \end{aligned}$$

Naturality

Suppose

$$\mathrm{hom}_{\mathcal{C}}(FX, Y) \cong \mathrm{hom}_{\mathcal{D}}(X, GY), \quad X \in \mathrm{Ob} \mathcal{C}, Y \in \mathrm{Ob} \mathcal{D}.$$

We say these isomorphisms are **natural** if the diagram

$$\begin{array}{ccc} \mathrm{hom}_{\mathcal{C}}(FX_2, Y_1) & \xrightarrow{\cong} & \mathrm{hom}_{\mathcal{D}}(X_2, GY_1) \\ \mu_{F(f),g} \downarrow & & \downarrow \mu_{f,G(g)} \\ \mathrm{hom}_{\mathcal{C}}(FX_1, Y_2) & \xrightarrow{\cong} & \mathrm{hom}_{\mathcal{D}}(X_1, GY_2) \end{array}$$

commutes for all

$$f \in \mathrm{hom}_{\mathcal{C}}(X_1, X_2), \quad g \in \mathrm{hom}_{\mathcal{D}}(Y_1, Y_2).$$

Adjoint functors: final definition

Definition (Adjoint functors)

A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is **left adjoint** to a functor $G: \mathcal{D} \rightarrow \mathcal{C}$ (and G is **right adjoint** to F) if we have **natural** isomorphisms

$$\mathrm{hom}_{\mathcal{D}}(FX, Y) \cong \mathrm{hom}_{\mathcal{C}}(X, GY), \quad X \in \mathrm{Ob}\mathcal{C}, Y \in \mathrm{Ob}\mathcal{D}.$$

Now that we've formulated the correct notion of adjoint functors, there is an obvious question:

Question: Should we care? Are there nice/familiar examples of adjoint functors?

Answer: YES!

Application: linear algebra

Let **Vect** denote the category of complex vector spaces.

Recall that we have the **forgetful functor**

$$\mathbf{forget} : \mathbf{Vect} \rightarrow \mathbf{Set}.$$

Let's try to find a **left adjoint** to **forget**.

So we want a functor $F : \mathbf{Set} \rightarrow \mathbf{Vect}$ such that we have natural isomorphisms

$$\mathrm{hom}_{\mathbf{Vect}}(FX, Y) \cong \mathrm{hom}_{\mathbf{Set}}(X, \mathbf{forget}(Y)),$$

for X a set and Y a vector space.

So, if X is a set and Y is a vector space, we want to associate to every **set** map $f : X \rightarrow Y$ a **vector space** map $Ff : FX \rightarrow Y$.

Application: linear algebra

Recall from linear algebra

Suppose V and W are vector spaces. If

- $v_i, i \in I$, is a basis of V , and
- $w_i, i \in I$, are **arbitrary** elements of W ,

then there is a **unique** linear map

$$A: V \rightarrow W \quad \text{such that} \quad Av_i = w_i.$$

Put another way: If

- B is a basis of V , and
- $f: B \rightarrow W$ is a **map of sets**,

then there is a unique way to extend f to a **linear map**

$$A: V \rightarrow W \quad \text{such that} \quad Av = fv \quad \forall v \in B.$$

Application: linear algebra

Recall: We want a functor $F: \mathbf{Set} \rightarrow \mathbf{Vect}$ such that we have natural isomorphisms

$$\mathrm{hom}_{\mathbf{Vect}}(FX, Y) \cong \mathrm{hom}_{\mathbf{Set}}(X, \mathbf{forget}(Y)),$$

for X a set and Y a vector space.

So FX should be a **vector space with basis X** .

Free vector spaces

Question: Given a set X , how do we form a vector space with basis X ?

For a set X , let

$$FX = \{f: X \rightarrow \mathbb{C} : f(x) = 0 \text{ for all but finitely many } x \in X\}.$$

For $x \in X$, we have the **Dirac delta function**:

$$\delta_x: X \rightarrow \mathbb{C}, \quad \delta_x(y) = \begin{cases} 1 & \text{if } y = x, \\ 0 & \text{if } y \neq x. \end{cases}$$

Exercise: The δ_x , $x \in X$, are a basis for FX .

So if we identify x with δ_x , then FX is a vector space with basis X .

FX is called the **free vector space on the set X** .

Example of adjoint functors in linear algebra

To every set map $f: X \rightarrow Y$, we can associate a natural linear map $Ff: FX \rightarrow FY$, and show that we have a functor **Set** \rightarrow **Vect**.

Theorem

The **free vector space functor**

$$F: \mathbf{Set} \rightarrow \mathbf{Vect}$$

is left adjoint to the **forgetful functor**

$$\mathbf{forget}: \mathbf{Vect} \rightarrow \mathbf{Set}.$$

We've done most of the work in proving this. One just needs to verify that the isomorphisms

$$\mathrm{hom}_{\mathbf{Vect}}(FX, Y) \cong \mathrm{hom}_{\mathbf{Set}}(X, \mathbf{forget}(Y)),$$

are natural (exercise).

Application: group theory

Let **Group** be the category of groups.

We have a forgetful functor

$$\mathbf{forget}: \mathbf{Group} \rightarrow \mathbf{Set}.$$

Just as for vector spaces, its left adjoint is the **free group functor**.

Suppose X is a set. The **free group** on X is the group FX with elements that are **words** in the alphabet X and formal inverses of elements of X :

$$a_1^{\pm 1} a_2^{\pm 1} \cdots a_n^{\pm 1}, \quad a_1, a_2, \dots, a_n \in X.$$

Multiplication is given by concatenation and cancelling inverses: e.g.

$$(aba^{-1}bc^{-1}ab)(b^{-1}ac) = aba^{-1}bc^{-1}aac.$$

Application: group theory

If $f: X \rightarrow Y$ is a map of sets, then we have the corresponding map of free groups:

$$FX \rightarrow FY, \quad a_1^{\pm 1} a_2^{\pm 1} \cdots a_n^{\pm 1} \mapsto f(a_1)^{\pm 1} f(a_2)^{\pm 1} \cdots f(a_n)^{\pm 1}.$$

So we have a **free group functor**

$$F: \mathbf{Set} \rightarrow \mathbf{Group}.$$

Theorem

The **free group functor**

$$F: \mathbf{Set} \rightarrow \mathbf{Group}$$

is left adjoint to the **forgetful functor**

$$\mathbf{forget}: \mathbf{Group} \rightarrow \mathbf{Set}.$$

Note: Can replace **Group** by the category of abelian groups to get the **free abelian group** on a set.

Application: adjoining an identity to a ring

Let **Ring** be the category of rings.

Let **Rng** be the category of **general rings** (rings without the multiplicative identity axiom).

The left adjoint to the **forgetful functor**

$$\mathbf{forget}: \mathbf{Ring} \rightarrow \mathbf{Rng}$$

is a functor that formally adjoins an identity to a ring. It maps a general ring R to $R \times \mathbb{Z}$ with multiplication determined by

$$\begin{aligned}(r, 0)(0, 1) &= (r, 0) = (0, 1)(r, 0), \\ (r, 0)(s, 0) &= (rs, 0), \\ (0, 1)(0, 1) &= (0, 1).\end{aligned}$$

This is a ring with identity $(0, 1)$.

Application: abelianization

Consider the inclusion functor

category of abelian groups \rightarrow **Group**.

It has a left adjoint called **abelianization**, which assigns to every group G the quotient group

$$G^{\text{ab}} = G/[G, G],$$

where $[G, G]$ is the subgroup of G generated by

$$g^{-1}h^{-1}gh, \quad g, h \in G.$$

A functor with both a left and right adjoint

Let **Top** be the category of topological spaces.

Consider the **forgetful functor**

$$\mathbf{forget}: \mathbf{Top} \rightarrow \mathbf{Set}.$$

Left adjoint: The functor taking a set to that same set with the **discrete topology** is left adjoint to **forget**.

Right adjoint: The functor taking a set to that same set with the **trivial topology** is right adjoint to **forget**.

Conclusion

Adjoint functors are a natural concept in mathematics that lead to many useful constructions:

- free vector spaces
- free (abelian) groups
- abelianization
- discrete/trivial topology

The idea of adjoint functors can be generalized further, to the setting of

- monoidal categories,
- 2-categories.

This is related to **higher representation theory**.